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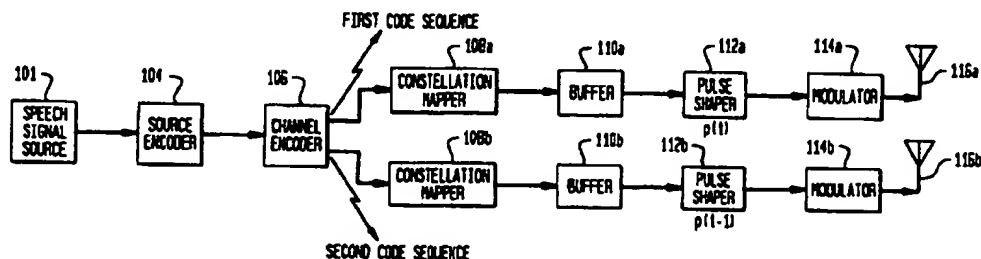
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(54) Title: METHOD AND APPARATUS FOR DATA TRANSMISSION USING MULTIPLE TRANSMIT ANTENNAS



## (57) Abstract

A method and apparatus for increasing the data rate and providing antenna diversity using multiple transmit antennas is disclosed. A set of bits of a digital signal are used to generate a codeword (106). Codewords are provided according to a channel code. Delay elements (110) may be provided in antenna output channels, or, with suitable code construction, delay may be omitted.  $n$  signals representing  $n$  symbols of a codeword are transmitted with  $n$  different transmit antennas. At the receiver MLSE (310) or other decoding is used to decode the noisy received sequence. The parallel transmission and channel coding enables an increase in the data rate over previous techniques, and recovery even under fading conditions. The channel coding may be concatenated with error correction codes under appropriate conditions.

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**METHOD AND APPARATUS FOR DATA TRANSMISSION  
USING MULTIPLE TRANSMIT ANTENNAS**

REFERENCE TO RELATED APPLICATIONS

This application claims priority from U.S.  
5 Provisional Application Serial Nos. 60/017,046 filed  
April 26, 1996 and 60/030,572 filed November 7, 1996.

BACKGROUND OF THE INVENTION

1. Field of Invention

The present invention relates generally to the  
10 field of communications systems, and particularly to the  
field of wireless communications, such as cellular radio.

2. Description of Related Art

Antenna diversity is a technique used in  
communication systems, including mobile cellular radio, to  
15 reduce the effects of multi-path distortion fading.  
Antenna diversity may be obtained by providing a receiver  
with two or more ( $n \geq 2$ ) antennas. These  $n$  antennas, when  
properly positioned, imply  $n$  channels which suffer fading  
in different manners. When one channel is in deep fade -  
20 that is, suffering severe amplitude and phase loss due to  
the destructive effects of multi-path interference,  
another of these channels is unlikely to be suffering from  
the same effect simultaneously. The redundancy provided by  
these independent channels enables a receiver to often  
25 avoid the detrimental effects of fading.

Alternatively, antenna diversity benefit can be  
provided to a mobile receiver by providing multiple  
transmitting antennas at a base or transmitting station,  
rather than at the receiver. The receiver can therefore  
30 use a single antenna, saving cost and complexity at that  
side of the transmission chain.

Multiple transmit antennas can be provided at the  
base station in a variety of ways. A schematic diagram of  
certain possible known techniques is illustrated in Figure  
35 1. Perhaps most simply, as schematically illustrated in  
Figure 1(a) two antennas can be provided at the output

stage, and the information <sup>2</sup> signal  $d_k$  can be switched between two matched antenna elements, without overlap in time or frequency. Of course this has the drawback that the transmitter requires feedback from the receiver about the channels corresponding to each transmit antenna. This scheme does not perform well when the channel is rapidly changing.

In a variant described in U.S. Patent No. 5,479,448 and schematically illustrated in Figure 1(b), the above mentioned drawbacks of switch diversity are removed by using a channel code to provide diversity benefit. Maximum diversity is upper-bounded by the number of antenna elements at the base station, and is equal to the minimum Hamming distance of the channel code used, provided that the receiver is equipped with one antenna. The system described in that patent is applicable to both FDD (frequency division duplex) and TDD (time division duplex)-based systems.

Illustrative embodiments of the system of U.S. Patent No. 5,479,448 comprise a base station which employs a channel code of length  $n \geq 2$  symbols ( $n$  being the number of antennas used by the transmitter), and a minimum Hamming distance  $2 \leq d_{\min} \leq n$ . This channel code is used to encode a group of  $k$  information bits. The  $n$  antennas of the base station transmitter are separated by a few wavelengths, as is conventional to provide the diversity reception with the  $n$  antennas. The channel code symbol  $c_i$  is transmitted with the  $i^{\text{th}}$  antenna to represent these  $k$  bits. At a receiver, a conventional maximum likelihood channel code decoder provides a diversity advantage of  $d_{\min}$ .

In the preferred embodiment of U.S. Patent No. 5,479,448, the transmitted signals from different antennas are separated in time. This results in data rate reduction, sacrificing bandwidth. The reduction in data

rate is equal to the number of antennas (or length of the code).

Transmit bandwidth can be improved over the diversity arrangement of Figure 1(b), by splitting the information signal into two paths to the two antennas, the second of which has a delay element or tap as disclosed in A. Wittneben, "Base Station Modulation Diversity for Digital SIMULCAST," 41<sup>st</sup> IEEE Vehicular Technology Society Conference Proceedings, pp. 848-853 and shown in Figure 1(c). The signal appearing at antenna B at any given instant of time is therefore the same signal as appeared at antenna A the preceding instant of time. The two signals are transmitted simultaneously, reconstructed at the receiving station, and processed to isolate the desired information signal.

#### SUMMARY OF THE INVENTION

The invention improving on these and other communication techniques in one aspect relates to a system and method for data transmission using multiple transmit antennas.

The invention in one aspect relates to a system and method for data transmission which increases effective utilization of available channel bandwidth, without great increases in transmitter or receiver complexity or cost.

The invention in another aspect relates to a system and method for data transmission which utilizes channel-codes to transmit data, reducing the chance of error and increasing reception robustness.

The invention in another aspect relates to a system and method for data transmission which can include concatenated error correcting codes, even further increasing BER and other transmission performance.

The invention in another aspect relates to a system and method for data transmission which can include multilevel coding, and decreases decoding complexity.

The invention in another aspect relates to a system and method for data transmission which preserves

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diversity benefit from multiple antenna arrangements, under a wide range of conditions.

In the present invention, among other advantages the time separation described in U.S. Patent No. 5,749,448 is removed, and coded data is transmitted in parallel, simultaneously from different transmit antennas, with or without delay. Increased data rate as well as diversity are achieved.

By way of comparison, the codes described in U.S. Patent No. 5,749,448 (col. 6, lines 21-29; col. 7, lines 35-44 and 63-67; col. 8, lines 1-16) provide a diversity 2 using 2 transmit antennas and 1 receive antenna. The bandwidth efficiencies for these disclosed codes are 1 bit/symbol, 1.5 bits/symbol and 2 bits/symbol respectively.

Using the present invention as described below, applying the same codes but a new transmission arrangement, the bandwidth efficiency doubles to 2, 3 and 4 bits/symbol respectively. Moreover, in another embodiment of the present invention when coding is done taking into account diversity and other criteria, no delay element on the antenna line is necessary to implement the invention and further coding gain is obtained.

#### BRIEF DESCRIPTION OF THE DRAWINGS

Figures 1(a), 1(b) and 1(c) illustrate a schematic diagram of certain prior approaches to multiple transmit antennas at base stations;

Figures 2(a) and 2(b) illustrate a schematic block diagram of first and second embodiments of multiple transmit antenna base stations, according to the invention;

Figure 3 illustrates a schematic block diagram of a wireless communication system constructed according to the illustrative first embodiment of the invention;

Figure 4 illustrates signal constellations used in implementations of the invention;

Figure 5 illustrates a schematic block diagram of a receiver constructed in conjunction with the first illustrative embodiment of the invention;

5 Figure 6 illustrates a schematic block diagram of decoding circuitry used in the receiver constructed according to the first illustrative embodiment of the invention;

10 Figure 7 illustrates a schematic block diagram of a receiver like that shown in Figure 5, but adapted to use two antenna elements;

Figure 8 illustrates a schematic block diagram of a wireless communication system constructed according to a second illustrative embodiment of the invention;

15 Figure 9 illustrates a 4-PSK code, used in implementation of the second illustrative embodiment of the invention;

20 Figure 10 illustrates a schematic block diagram of decoding circuitry used in a receiver constructed according to the second illustrative embodiment of the invention;

Figure 11 illustrates an 8-PSK code, used in implementation of the second illustrative embodiment of the invention;

25 Figure 12 illustrates a 4-PSK code with 8 and 16 states, used in implementation of the second illustrative embodiment of the invention;

Figure 13 illustrates a 4-PSK code with 32 states, used in implementation of the second illustrative embodiment of the invention;

30 Figure 14 illustrates a 2-Space-Time QAM code with 16 and 16 states, used in implementation of the second illustrative embodiment of the invention;

35 Figure 15 illustrates data demonstrating transmission performance of transmission according to the second illustrative embodiment of the invention; and

6  
Figure 16 illustrates a time slot structure related to channel probing techniques used in connection with the invention;

5 Figure 17 illustrates a schematic diagram of a transmitter that employs space-time coding with 2 transmit antennas;

Figure 18 illustrates a schematic diagram of the receiver with space-time vector Viterbi decoder;

10 Figure 19 illustrates the frame-error-rate performance of the basic modem structure;

Figure 20 shows the estimated distribution of the number of symbol errors per frame at Doppler frequency 170 Hz;

15 Figure 21 illustrates a schematic diagram for the transmitter with concatenated space-time coding according to a third illustrative embodiment of the invention;

20 Figure 22 illustrates a schematic diagram for the receiver with space-time vector Viterbi decoder concatenated with a Reed-Solomon decoder according to the third illustrative embodiment;

Figure 23 illustrates the performance of the concatenated space-time code of the third illustrative embodiment of the invention;

25 Figure 24 describes set partitioning of a 16 QAM constellation to be used in an example of multi-level space-time codes according to the fourth illustrative embodiment of the invention;

Figure 25 describes example of encoders for different levels of multi-level space-time code;

30 Figure 26 describes an equivalent space-time code for an example of a multi-level space-time code constructed according to the fourth illustrative embodiment of the invention; and

35 Figures 27(a) and 27(b) respectively illustrate smart greedy codes constructed using the BPSK and 4-PSK constellations, according to a fifth illustrative embodiment of the invention.



DETAILED DESCRIPTION OF PREFERRED EMBODIMENTSA. Incorporation by Reference

Various concepts of digital signal processing referred to in this application are well known in, for example, the digital communication and other arts, and thus they need not be described in detail herein. These concepts include, without limitation, combined modulation and coding, and maximum-likelihood decoding. These concepts are described for instance in U.S. Patent No. 4,457,004, issued June 26, 1984 to A. Gersho et al.; U.S. Patent No. 4,489,418, issued December 18, 1984 to J.E. Mazo; U.S. Patent No. 4,520,490, issued May 28, 1985 to L. Wei; U.S. Patent No. 4,597,090, issued June 24, 1986 to G.D. Forney, Jr.; U.S. Patent No. 5,029,185 issued July 2, 1991 to L. Wei; in A. Wittneben, "Base Station Modulation Diversity for Digital SIMULCAST," 41<sup>st</sup> IEEE Vehicular Technology Society Conference Proceedings, pp. 848-853; and U.S. Patent No. 5,479,448 to Seshadri, all of which are incorporated by reference.

B. Illustrative Hardware used in Embodiments

For clarity of explanation, illustrative embodiments of the present invention are presented as comprising individual functional blocks. As known in the art, the functions these blocks represent may be provided through the use of either shared or dedicated hardware (processors), including, but not limited to, hardware capable of executing software. Illustrative embodiments may comprise digital signal processor (DSP) hardware, and software performing the operations discussed below. Very large scale integration (VLSI) hardware embodiments of the present invention, as well as hybrid DSP/VLSI embodiments, may also be constructed.

C. Introduction to Illustrative Embodiments

The central idea of conventional antenna diversity reception is that with high probability, a signal received at different antennas undergoes fading at different moments in time. Thus, a receiver can combine or select

different receive signals to reconstruct the transmitted signal with little distortion.

The present invention provides the benefit of diversity by taking advantage of multiple antennas at a transmitter, with or without delay. A first illustrative embodiment shown in Figures 2(a) and 3 maps the information sequence of length  $M_1$  to a two code sequence of length  $M_2$ . In particular every group of  $k$  input bits (assume  $k$  divides  $M_1$ ) are mapped to first and second code symbols. The two code symbols are used to form two code sequences where each sequence is of length  $M_1/k=M_2$  where the first code sequence is comprised of the first code symbol while the second one is comprised of the second code symbol. These two code sequences are then used to phase modulate a carrier using conventional phase shift keying, as is well known in the art, and in that process two modulated signals are generated. Alternatively, quadrature amplitude modulation, or any other modulation scheme can be used.

The two modulated signals are then transmitted using two transmit antennas. In the first illustrative embodiment, a timing offset of one symbol interval (delay element or tap, of period  $T$ ) is introduced between the two signals. The receiver receives a sum of faded versions of the transmitted signals from the two antennas, perturbed by noise. In the second illustrative embodiment, the use of a delay in one of the antenna channels is eliminated.

Because the two coded signals are transmitted simultaneously, no bandwidth penalty is incurred. However intersymbol interference is created which is resolved at the receiver using maximum likelihood sequence detection or other techniques that are known in the art. As noted, the introduction of delay to provide diversity is known in the art. However the use of coding as an integral part of the delay diversity arrangement is not known, nor is elimination of any delay element using codes which adhere to diversity and other criteria.

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Prior to proceeding with a description of illustrative embodiments of the present invention, concepts related to a channel model for the first illustrative embodiment and embodiment error performance will be presented.

D. Channel Model Transmission Format:  
Analysis in First Illustrative Embodiment

The overall transmission environment in which the first illustrative embodiment of the invention operates may be viewed as comprising  $n$  distinct channels, each illustratively undergoing independent slow (static) Rayleigh fading (it should be understood that the principles of the present invention are applicable to other classes of fading channels as well). The impulse response for the  $i^{\text{th}}$  channel is given by

$$h_i(t) = \alpha_i \delta(t) e^{j\omega_0 t}, \quad 1 \leq i \leq N \quad (1)$$

where  $\omega_0$  is the angular carrier frequency and  $\alpha_i$  is the static complex fade value whose phase is a random variable that is uniformly distributed over  $[-\pi, \pi)$ , and whose magnitude is Rayleigh distributed with

$$P(|\alpha_i|) = 2|\alpha_i| e^{-|\alpha_i|^2}, \quad |\alpha_i| \geq 0 \quad (2)$$

The information sequence  $I$  is grouped into sub-sequences of  $k$  information bits,

$$I = \left( \underbrace{I_0^1, I_1^1, I_2^1, \dots, I_{k-1}^1}_{\text{1st sub-sequence}}, \underbrace{I_0^2, \dots, I_{k-1}^2}_{\text{2nd sub-sequence}}, \dots \right)$$

where the superscript is the sub-sequence number. Each sub-sequence is mapped into  $n$  channel symbols of the channel constellation using a channel code. Some of the illustrative signal constellations are shown in Figure 4. The signal constellation mapped code sequence is

10

$$c = \left( \begin{array}{cc} \underbrace{c_0^1, \dots, c_{n-1}^1}_{\text{code sequence for 1st sub-sequence}} & \underbrace{c_0^2, \dots, c_{n-1}^2}_{\text{code sequence for 2nd sub-sequence}} \dots \end{array} \right)$$

Hence each element  $c_i^j$  is a point belonging to a signal constellation. The code sequence is arranged in a matrix as shown below

$$\begin{bmatrix} c_0^1 & c_0^2 & c_0^3 & \dots \\ c_1^1 & c_1^2 & c_1^3 & \\ \vdots & & & \\ c_{n-1}^1 & c_{n-1}^2 & c_{n-1}^3 & \dots \end{bmatrix}$$

The first row of the matrix is pulse shaped using square-root Nyquist filter  $p(t)$ , modulated and transmitted using antenna 1. The second row of the matrix is pulse shaped using square-root Nyquist filter  $p(t-T)$  ( $p(t)$  delayed by one symbol interval). The  $i^{\text{th}}$  row of the matrix is transmitted using square root Nyquist filter  $p(t-(i-1)T)$  ( $p(t)$  delayed by  $(i-1)$  symbol intervals). At the receiver, the received signal, following demodulation, receiver filtering and sampling as is well known in the art, is given by

$$r_i = \alpha_0 c_0^i + \alpha_1 c_1^{i-1} + \alpha_2 c_2^{i-2} + \dots + \alpha_{n-j} c_{n-j}^{i-(n-1)} + \eta_i$$

where  $\eta_i$  is the extraneous noise which is modeled as additive white Gaussian.

Decoding is done in a conventional manner using maximum likelihood decoding techniques or suboptimum variants thereof, which are well known in the art.

#### E. First Illustrative Embodiment

Figure 3 presents an illustrative apparatus of a digital wireless communication system transmitter according to a first illustrative embodiment of the

present invention. The transmitter receives an analog speech signal from speech signal source 101, and processes this signal for transmission on antennas 116a,b. The transmitter comprises a source encoder 104, a channel encoder 106, constellation mappers 108a,b, temporary storage buffers 110a,b, pulse shapers 112a and b, and modulators 114a,b. Power amplification associated with the transmission of radio signals has been omitted from Figure 3 for clarity.

The speech signal source 101 provides an analog speech signal to be encoded and transmitted for instance to a mobile receiver. This speech signal is converted to a digital signal by conventional analog-to-digital conversion by source encoder 104. Source encoder 104 provides a digital signal representative of the analog speech signal as output to channel encoder 106. Source encoder 104 may be realized with any of the conventional speech encoders.

The channel encoder 106 receives the PCM (Pulse Code Modulated) digital signal comprising a plurality of bits from the source encoder 104. Channel encoder 106 codes the PCM digital signal using a conventional channel code. Any channel code may be employed for this purpose, as long as it is appropriately constructed.

The code constructed for the first illustrative embodiment of the present invention assumes that the number of antennas at the base station is two. The following illustrative code of length  $n = 2$  complex symbols ( $2 \text{ symbols} \times 2 \text{ components (in-phase and quadrature) per symbol equals 4 dimensions (4-D)}$ ), has a minimum Hamming distance  $d_{min}=2$ .

Channel Code		
Information Bits	Symbol 1	Symbol 2
00	0	0
01	1	2
11	2	1
10	3	3

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Using this code, encoder 106 codes two information bits at a time to generate one of four codewords. Each generated codeword comprises two symbols (see columns labeled Symbol 1 and Symbol 2, above). Each symbol belongs to the 4-PSK constellation presented in Figure 4(a). Thus, a coding rate of one information bit per code symbol is provided by this code. Symbol 1 is transmitted with antenna 116a and symbol 2 with antenna 116b, as discussed below.

The first symbol of each codeword generated by encoder 106 is provided as input to constellation mapper 108a, and the second symbol of the codeword is provided to mapper 108b.

Constellation mappers 108a, b produce a complex valued output corresponding to a symbol received from encoder 106. The real part of this output determines an in-phase component of a modulated signal transmitted at antennas 116a,b. Similarly, the imaginary part of the output determines a quadrature component of the modulated signal. The constellation mapper 108a,b are conventional mappers known in the art. They may be realized as a look-up table or as a straightforward combination of logic elements. Mappers 108a,b operate on the first and second symbol of each received codeword, respectively, and provide complex valued output to buffers 110a and b.

Buffers 110a and b provide temporary storage for the complex values received from mappers 108a, b, and illustratively store 100 of such values. The complex entries in buffer 110a are pulse shaped using conventional square-root Nyquist transmit filter (see 112a) while those in buffer 110b are pulse shaped using the same square-root Nyquist transmit filter but whose impulse response is delayed by one symbol interval (see 112b). The pulse shaped outputs are then modulated by modulators 114a and 114b and transmitted using antennas 116a and 116b.

Additional filtering and power <sup>13</sup> amplification stages are not shown for clarity.

F. Further Channel Codes for  
First Illustrative  
Embodiment

The first embodiment described above may employ other channel codes than the one first developed, to enhance coding efficiency. For example, the following code length 2,  $d_{\min}=2$ , is formed from an 8-PSK constellation shown in Figure 4(b). This code has efficiency of 3 bits/symbol:

Information Data	Symbol 1	Symbol 2
000	0	0
001	1	5
011	2	2
111	3	7
100	4	4
101	5	1
110	6	6
111	7	3

A distinct pair of codewords differ in at least two positions.

In another coding implementation, a coding efficiency of 4.0 bits/symbol is provided. In order to achieve  $d_{\min}=2$  and stay within the constraint that the block length of the code equal two, it is necessary to have at least 16 codewords. Hence, 16-PSK (see Figure 4(c)) is the smallest constellation with which a diversity benefit of 2 can be provided. The 4D-16 PSK code is shown below:

Information Data	Symbol 1	Symbol 2
0000	0	0
0001	2	2
0010	4	4
0011	6	6
0100	8	8
0101	10	10
0110	12	12
0111	14	14
1000	1	7
1001	3	9
1010	5	11
1011	7	13

Information Data	Symbol 1	Symbol 2
1100	9	15
1101	11	1
1110	13	3
1111	15	5

G. An Illustrative Decoder for Embodiments

Figure 5 presents an illustrative receiver 300 according to the foregoing first illustrative embodiment of the present invention. Receiver 300 receives transmitted signals from antenna 301, and produces analog speech as output. Receiver 300 comprises an RF-to-baseband front end 305, receive buffer 307, channel decoder 310, and speech decoder 320.

The RF-to-baseband front end 305 provides conventional demodulated output (i.e., received symbols) to the receive buffers 307. Front end 305 includes, e.g., conventional RF to IF conversion, receive filtering, and timing and carrier recovery circuits.

Receive buffer 307 store received symbols from front end 305. Buffer 307 analogous to buffers 110a, b of the illustrative transmitter described in Section D and present in Figure 3 except that since the receiver receives a superposition of data in buffers 110a, b only one buffer is needed. Channel decoder 210 receives the demodulated symbol output from buffer 307, and provides decoded information bits to speech decoder 320. The illustrative decoder 310 operates in accordance with the flow diagram presented in Figure 6.

As shown in Figure 6, symbols from receive buffer 307 are used in computing distances with all possible valid codewords stored in memories 311a, b. For example the first codeword from buffer 311a taken together with the first codeword from 311b, but delayed by one unit symbol interval are linearly combined with channel gains  $\alpha_1$  and  $\alpha_2$  respectively. The distance between this combined output and the received symbols in buffer 307 is computed. This is done for every codeword in buffers 311a and 311b



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(see 312). The legal codeword pair is the one which most closely match the received sequence (see 313). The decoded codeword pair is then mapped to a string of bits which comprises coded information (see 314). This exhaustive search can be implemented efficiently using the Viterbi algorithm or variants thereof, known to persons skilled in the art.

Speech decoder 320 is a conventional device providing a mapping of digital speech information to analog speech. Decoder 320 provides an inverse operation to source encoder 104 discussed above with respect to Figure 5.

In light of the discussion above, it is to be understood that the diversity benefit of the present invention using one antenna may be enhanced by use of multiple receive antennas. This advantage may be realized by combination of a front end and receive buffer for each receiver antenna.

Figure 7 presents an illustrative decoder in accordance with this enhancement for two receiving antennas 301a, b. As shown in the figure, received symbols from the first and second buffers associated with each antenna are provided directly to channel decoder. These are processed in a manner similar to the one described above by the decoder and a decision on the transmitted signal is made.

H. Second Illustrative Embodiment: Introduction

In the present invention, the foregoing first illustrative embodiment of the invention and its coding implementations rely upon coding technique and a delay element in the antenna transmission line, to preserve diversity and achieve additional coding gain over the simpler known delay diversity schemes. However, that illustrative embodiment can be further improved by removing the restriction that delays be introduced between different coded streams.

In particular, in <sup>16</sup>th second illustrative embodiment of the invention, the inventors derive criteria for maximizing the performance when  $n$  transmit antennas are used to transmit  $n$  parallel data streams that are created by encoding information data with a channel code. In particular, it is shown that the code's performance is determined by the rank and determinant of certain matrices. These matrices in turn are constructed from codewords of the given channel code. These matrix based criteria are used to design channel codes for high data rate wireless communications. These codes are called space-time codes, and are easy to encode because they have a trellis structure. These codes can be easily decoded using maximum likelihood sequence criterion. Examples of 4PSK, 8PSK and 16QAM based codes are given that have been constructed for operation with 2 and 4 transmit antennas. Performance results are shown to verify the performance.

I. Channel Model Transmission Format:  
Analysis for Second Illustrative  
Embodiment

The overall transmission channel in which the second illustrative embodiment and its coding implementation operates may be viewed as comprising  $n$  distinct channels, each illustratively undergoing independent slow (static) Rayleigh or Rician fading (it should again be understood that the principles of the present invention and this embodiment are applicable to other classes or fading channels as well), having impulse response, fade and other characteristics generally as described above for the first illustrative embodiment.

J. Second Illustrative Embodiment

Figure 8 presents a communication system constructed according to the second illustrative of the present invention. The system shown is generally similar to that of the first illustrative embodiment shown in Fig. 3, and elements in common with the previous embodiment are labeled with similar numbers, including signal source 101,

antennas 116a,b, encoder 104 and chann 1 encoder 106, and  
constellation mappers 108a,b. It may be noted that pulse  
shaper 112b' in the second illustrative embodiment is not  
constructed to apply a delay of T, but is the same as  
5 pulse shaper 112a'.

The channel encoder 106 receives the PCM digital  
signal comprising a plurality of bits from the source  
encoder 104. Channel encoder 106 codes the PCM digital  
signal using a channel code that has been constructed to  
10 meet the design criteria elucidated below.

The code constructed for the second illustrative  
embodiment assumes that the number of antennas at the base  
station is two. The 4-PSK trellis code with a  
transmission rate of 2 bits/sec/Hz is provided for  
15 illustrative purposes in Fig. 9. Using this code, encoder  
106 codes two information bits at a time to generate the  
label of a branch in the trellis diagram. The branch  
depends on the state of the encoder and the input data and  
determines the new state of the encoder as well. For  
20 example, suppose that the encoder is in state 3 of Fig. 9.  
Then upon input bits, 00, 01, 10, and 11, the respective  
branch labels are respectively 30, 31, 32, and 33. The  
new state of the encoder is then respectively 0, 1, 2, and  
3. Each branch label comprises two symbols (see branch  
25 labels, above). Each symbol belongs to the 4-PSK  
constellation presented in Fig. 4(a). Thus for instance  
corresponding to output 31, phase values  $3\pi/2$  and  $\pi/2$   
radians are used to phase modulate the carrier.  
Therefore, a coding rate of two information bits per  
30 channel used is provided by this code. Symbol 1 is  
transmitted with antenna 116a and symbol 2 with antenna  
116b, as discussed below.

The first symbol of each codeword generated by  
encoder 106 is provided as input to constellation mapper  
35 108a, and the second symbol of the codeword is provided  
to mapper 108b, generally as discussed above for the  
first illustrative embodiment.

18

K. Further Illustrative Channel  
Codes in Second Illustrative  
Embodiment

The second illustrative embodiment described  
5 above may employ other channel codes to enhance coding  
efficiency. These codes are designed according to a  
performance criteria computed later in the sequel. For  
illustration, examples are provided. One can improve on  
the performance of these codes by constructing encoders  
10 with more states. The inventors have designed codes  
(using the criteria established) with different numbers  
of states. Simulation results for the case of 4-PSK and  
8-PSK are included demonstrating that the performance of  
these codes for two and one receive antenna is excellent.

15 L. Decoding in Second Illustrative Embodiment

The second illustrative embodiment makes use of  
receiver 300 and related decoder circuitry illustrated in  
Fig. 10, generally similar to that shown in Fig. 5  
described for the first illustrative embodiment. As  
20 illustrated in Fig. 10, the circuitry constructed to  
receive symbols from buffer 307 is adapted to account for  
the non-delayed coding of the second embodiment. For  
instance, since no delay is applied, the delay element  
315 shown in Fig. 6 is not incorporated when decoding  
25 according to the second illustrative embodiment.

M. Performance Criteria for  
Second Illustrative  
Embodiment

In this section, performance criteria for the  
30 design of the codes used in the second illustrative  
embodiment are established.

Consider a mobile communication system such that  
the base station is equipped with  $n$  antennas and the  
mobil unit is equipped with  $m$  antennas. Data is encoded  
35 by the encoder. The encoded data goes through a serial  
to parallel device and is divided into  $n$  streams of data.  
Each stream of data is used as the input to a pulse

shaper. The output of each shaper<sup>19</sup> is then modulated using a modulator. At each time the output of modulator  $i$  is a signal that is transmitted using transmit antenna (Tx antenna)  $i$  for  $1 \leq i \leq n$ .

5 It is again assumed that the  $n$  signals are transmitted simultaneously each from a different transmit antenna and that all these signals have the same transmission period  $T$ . The signal at each receive antenna is a noisy version of the superposition of the  
10 faded version of the  $n$  transmitted signals.

At the receiver, the demodulator makes a decision statistic based on the received signals at each receive antenna  $1 \leq j \leq m$ . Assuming that the transmitted symbol from the  $i$ -th antenna at transmission interval  $t$  is  $c_t^i$ ,  
15 and the receive word at time interval  $t$  at the receive antenna  $j$  is  $d_t^j$ , then

$$d_t^j = \sum_{i=1}^n \alpha_i^j c_t^i + \eta_t^j \quad (3)$$

The coefficients  $\alpha_i^j$  are first modeled as independent samples of a stationary complex Gaussian stochastic  
20 process with mean  $E\alpha_i^j = p_i^j + q_i^j j$  and variance 0.5 per dimension with  $K_i^j = |E\alpha_i^j|^2$ , where  $j = \sqrt{-1}$ . This is equivalent to the assumption that signals transmitted from different antennas undergo independent fades (The case when  $\alpha_i^j$  are dependent will be treated later). Also,  $\eta_t^j$   
25 are independent samples of a zero mean complex white Gaussian process with two sided power spectral density  $N_0/2$  per dimension. It is assumed that  $\alpha_i^j$  are constant during a frame and vary from one frame to another (flat fading).

30 The inventors have derived a design criterion for constructing codes under this transmission scenario. Mathematical background required and the notation used

20

for this task is first reviewed. Let  $\mathbf{x} = (x_1, x_2, \dots, x_k)$  and  $(y_1, y_2, \dots, y_k)$  be complex vectors in the  $k$  dimensional complex space  $C^k$ . The inner product  $\mathbf{x}$  and  $\mathbf{y}$  is given by

$$\mathbf{x} \cdot \mathbf{y} = \sum_{i=1}^k x_i \bar{y}_i, \text{ where } \bar{y}_i \text{ denotes the complex conjugate of } y_i.$$

5 For any matrix  $A$ , let  $A^*$  denote the Hermitian (transpose conjugate) of  $A$ .

From known linear algebra an  $n \times n$   $A$  is *Hermitian* if and only if  $A = A^*$ .  $A$  is *non-negative definite* if  $\mathbf{x}A\mathbf{x}^* \geq 0$  for any  $1 \times n$  complex vector  $\mathbf{x}$ . An  $n \times n$  matrix  $V$  is *unitary* if and only if  $VV^* = I$  where  $I$  is the identity matrix. A  $n \times 1$  matrix  $B$  is a *square root* of an  $n \times n$  matrix  $A$  if  $BB^* = A$ . The following results from linear algebra are also made use of.

- An eigenvector  $\mathbf{v}$  of an  $n \times n$  matrix  $A$  corresponding to eigenvalue  $\lambda$  is a  $1 \times n$  vector of unit Euclidean length such that  $\mathbf{v}A = \lambda\mathbf{v}$  for some complex number  $\lambda$ . The number of eigenvectors of  $A$  corresponding to the eigenvalue zero is  $n - r$ , where  $r$  is the rank of  $A$ .
- 20 • Any matrix  $A$  with a square root  $B$  is non-negative definite.
- For any non-negative definite Hermitian matrix  $A$ , there exists a lower triangular square matrix  $B$  such that  $BB^* = A$ .
- 25 • Given a Hermitian matrix  $A$ , the eigenvectors of  $A$  span  $C^n$ , the complex space of  $n$  dimensions and it is easy to construct an orthonormal basis of  $C^n$  consisting of eigenvectors of  $A$ .
- There exists a unitary matrix  $V$  and a real diagonal matrix  $D$  such that  $VAV^* = D$ . The rows of  $V$  are an orthonormal basis of  $C^n$  given by eigenvectors of  $A$ .
- 30 • The diagonal elements of  $D$  are the eigenvalues  $\lambda_i$ ,  $i = 1, 2, \dots, n$  of  $A$  counting multiplicities.

- The eigenvalues of a <sup>2</sup> Hermitian matrix are real.
- The eigenvalues of a non-negative definite Hermitian matrix are non-negative.

i. The Case of Independent Fade Coefficients

5 Assume that each element of signal constellation is contracted by a scale factor  $\sqrt{E_s}$ , chosen so that the average energy of the constellation element is 1. Thus the design criterion is not constellation dependent and applies equally to 4-PSK, 8-PSK and 16-QAM.

10 Consider the probability that the receiver decides erroneously in favor of a signal

$$\mathbf{e} = e_1^1 e_1^2 \cdots e_1^n e_2^1 e_2^2 \cdots e_2^n \cdots e_1^1 e_1^2 \cdots e_1^n$$

assuming that

$$\mathbf{c} = c_1^1 c_1^2 \cdots c_1^n c_2^1 c_2^2 \cdots c_2^n \cdots c_1^1 c_1^2 \cdots c_1^n$$

15 was transmitted.

Assuming ideal channel state information (CSI), the probability of transmitting  $\mathbf{c}$  and deciding in favor of  $\mathbf{e}$  at the decoder is well approximated by

$$P(\mathbf{c} \rightarrow \mathbf{e} | \alpha_i^j a_i^j, i=1,2,\dots,n, j=1,2,\dots,m) \leq \exp(-d^2(\mathbf{c}, \mathbf{e}) E_s / 4N_0) \quad (4)$$

20 where

$$d^2(\mathbf{c}, \mathbf{e}) = \sum_{j=1}^m \sum_{t=1}^l \left| \sum_{i=1}^n \alpha_i^j (c_t^i - e_t^i) \right|^2. \quad (5)$$

This is just the standard approximation to the Gaussian tail function.

Setting  $\Omega_j = (\alpha_1^j, \dots, \alpha_n^j)$ , (5) is rewritten as

25 
$$d^2(\mathbf{c}, \mathbf{e}) = \sum_{j=1}^m \Omega_j A(\mathbf{c}, \mathbf{e}) \Omega_j^*, \quad (6)$$

where the pq in element of  $A(\mathbf{c}, \mathbf{e})$  is  $A_{pq} = \mathbf{x}_p \cdot \mathbf{x}_q$  and  $\mathbf{x}_p = (c_1^p - e_1^p, c_2^p - e_2^p, \dots, c_1^p - e_1^p)$  for  $1 \leq p, q \leq n$ . Thus,

2.2

$$P(\mathbf{c} \rightarrow \mathbf{e} \mid \alpha_i^j, i=1,2,\dots,n, j=1,2,\dots,m) \leq \prod_{j=1}^m \exp(-\Omega_j A(\mathbf{c}, \mathbf{e}) \Omega_j^* E_j / 4N_0). \quad (7)$$

Since  $A(\mathbf{c}, \mathbf{e})$  is Hermitian, there exists a unitary matrix  $V$  and a real diagonal matrix  $D$  such that  $VA(\mathbf{c}, \mathbf{e})V^* = D$ . The rows of  $V$  are a complete orthonormal basis of  $\mathbb{C}^n$  given by eigenvectors of  $A$ . Furthermore, the diagonal elements of  $D$  are the eigenvalues  $\lambda_i, i=1,2,\dots,n$  of  $A$  counting multiplicities. The matrix

$$B(\mathbf{c}, \mathbf{e}) = \begin{pmatrix} e_1^1 - c_1^1 & e_2^1 - c_2^1 & \dots & \dots & e_n^1 - c_n^1 \\ e_1^2 - c_1^2 & e_2^2 - c_2^2 & \dots & \dots & e_n^2 - c_n^2 \\ e_1^3 - c_1^3 & e_2^3 - c_2^3 & \ddots & \vdots & e_n^3 - c_n^3 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ e_1^n - c_1^n & e_2^n - c_2^n & \dots & \dots & e_n^n - c_n^n \end{pmatrix} \quad (8)$$

is clearly a square root of  $A(\mathbf{c}, \mathbf{e})$ . Thus, the eigenvalues of  $A(\mathbf{c}, \mathbf{e})$  are non-negative real numbers.

Let  $\omega_j = \Omega_j V^*$  and  $\omega_j = (\beta_1^j, \dots, \beta_n^j)$ , then

$$\Omega_j A(\mathbf{c}, \mathbf{e}) \Omega_j^* = \sum_{i=1}^n \lambda_i |\beta_i^j|^2. \quad (9)$$

Next, recall that  $\alpha_i^j$  are i.i.d. samples of a complex Gaussian process with mean  $E\alpha_i^j$  with  $K_i^j = |E\alpha_i^j|^2$ . Let  $\mathbf{K}^j = (E\alpha_1^j, \dots, E\alpha_n^j)$ , and let  $\mathbf{v}_w$  denote the  $w$ -th row of  $V$ .

Since  $V$  is unitary,  $(\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n)$  is an orthonormal basis of  $\mathbb{C}^n$  and  $\beta_i^j$  are independent complex Gaussian random variables with variance 0.5 per dimension and mean  $\mathbf{K}^j \cdot \mathbf{v}_i$ . Let  $K_{i,j} = |E\beta_i^j|^2 = |\mathbf{K}^j \cdot \mathbf{v}_i|^2$ . Thus  $|\beta_i^j|$  are independent Rician distributions with pdf

$$p(|\beta_i^j|) = 2|\beta_i^j| \exp(-|\beta_i^j|^2 - K_{i,j}) I_0(2|\beta_i^j| \sqrt{K_{i,j}}),$$

for  $|\beta_i^j| \geq 0$ , where  $I_0(\cdot)$  is the zero-order modified Bessel function of the first kind.



Thus, to compute an <sup>23</sup> upper bound on the average probability of error, simply average  $\prod_{j=1}^m \exp(E_s / 4N_0) \sum_{i=1}^n \lambda_i |\beta_j^i|^2$  with respect to independent Rician distributions of  $|\beta_j^i|$  to arrive at

$$5 \quad P(\mathbf{c} \rightarrow \mathbf{e}) \leq \prod_{j=1}^m \left( \prod_{i=1}^n \frac{1}{1 + \frac{E_s}{4N_0} \lambda_i} \exp \left( - \frac{K_{i,j} \frac{E_s}{4N_0} \lambda_i}{1 + \frac{E_s}{4N_0} \lambda_i} \right) \right) \quad (10)$$

Some special cases are next examined.

**The Case of Rayleigh Fading:** In this case  $K_{i,j}^j = 0$  and as a fortiori  $K_{i,j} = 0$  for all  $i$  and  $j$ . Then the inequality (10) can be written as

$$10 \quad P(\mathbf{c} \rightarrow \mathbf{e}) \leq \left( \frac{1}{\prod_{i=1}^n (1 + \lambda_i E_s / 4N_0)} \right)^m. \quad (11)$$

Let  $r$  denote the rank of matrix  $A$ , then the kernel of  $A$  has dimension  $n-r$  and exactly  $n-r$  eigenvalues of  $A$  are zero. Say the nonzero eigenvalues of  $A$  are  $\lambda_1, \lambda_2, \dots, \lambda_r$ , then it follows from inequality (11) that

$$15 \quad P(\mathbf{c} \rightarrow \mathbf{e}) \leq \left( \prod_{i=1}^r \lambda_i \right)^{-m} (E_s / 4N_0)^{-rm}. \quad (12)$$

Thus a diversity of  $mr$  and a gain of  $(\lambda_1 \lambda_2 \dots \lambda_r)^{1/r}$  is achieved. Recall that  $\lambda_1 \lambda_2 \dots \lambda_r$  is the absolute value of the sum of determinants of all the principle  $r \times r$  cofactors of  $A$ . Moreover, the ranks of  $A(\mathbf{c}, \mathbf{e})$  and  $B(\mathbf{c}, \mathbf{e})$  are equal. Thus from the above analysis, the following design criterion are arrived at.

#### **Design Criteria For Rayleigh Space-Time Codes:**

- **The Rank Criterion:** In order to achieve the maximum diversity  $mn$ , the matrix  $B(\mathbf{c}, \mathbf{e})$  has to be full rank for any codewords  $\mathbf{c}$  and  $\mathbf{e}$ . If  $B(\mathbf{c}, \mathbf{e})$  has minimum rank  $r$  over the set of two tuples of

distinct codewords, <sup>24</sup> then a diversity of  $rm$  is achieved.

- **The Determinant Criterion:** Suppose that a diversity benefit of  $rm$  is our target. The minimum of  $r$ -th roots of the sum of determinants of all  $r \times r$  principle cofactors of  $A(\mathbf{c}, \mathbf{e}) = B(\mathbf{c}, \mathbf{e})B^*(\mathbf{c}, \mathbf{e})$  taken over all pairs of distinct codewords  $\mathbf{e}$  and  $\mathbf{c}$  corresponds to the coding gain, where  $r$  is the rank of  $A(\mathbf{c}, \mathbf{e})$ . Special attention in the design must be paid to this quantity for any codewords  $\mathbf{e}$  and  $\mathbf{c}$ . The design target is making this sum as large as possible. If a diversity of  $nm$  is the design target, then the minimum of the determinant of  $A(\mathbf{c}, \mathbf{e})$  taken over all pairs of distinct codewords  $\mathbf{e}$  and  $\mathbf{c}$  must be maximized.

At sufficiently high signal to noise ratios, one can approximate the right hand side of inequality (10) by

$$P(\mathbf{c} \rightarrow \mathbf{e}) \leq \left(\frac{E_s}{4N_0}\right)^{-rm} \left(\prod_{i=1}^r \lambda_i\right)^{-m} \left[ \prod_{j=1}^m \prod_{i=1}^r \exp(-K_{i,j}) \right]. \quad (14)$$

- Thus a diversity of  $rm$  and a gain of  $(\lambda_1 \lambda_2 \cdots \lambda_r)^{1/r} \left[ \prod_{j=1}^m \prod_{i=1}^r \exp(-K_{i,j}) \right]^{1/rm}$  is achieved. Thus, the following design criteria is valid for the Rician space-time codes for large signal to noise ratios.

**Design Criteria For The Rician Space-Time Codes:**

- **The Rank Criterion:** This criterion is the same as that given for the Rayleigh channel.
- **The Gain Criterion:** Let  $\Lambda(\mathbf{c}, \mathbf{e})$  denote the sum of all the determinants of  $r \times r$  principal co-factors of  $A(\mathbf{c}, \mathbf{e})$ , where  $r$  is the rank of  $A(\mathbf{c}, \mathbf{e})$ . The minimum of the products

$$\Lambda(\mathbf{c}, \mathbf{e})^{1/r} \left[ \prod_{j=1}^m \prod_{i=1}^r \exp(-K_{i,j}) \right]^{1/rm}$$

Taken over distinct codewords  $c$  and  $e$  have to be maximized.

- Note that it has been shown that, one could still use the gain criterion for the Rayleigh space-time codes as well, since the performance will be at least as good as the right side of inequality (11).

ii. The Case of Dependent Fade Coefficients:

Next, the case when the fade coefficients are dependent is studied. Only Rayleigh fading is considered, as the Rician case can be treated in a similar manner. To this end, consider the  $mn \times mn$  matrix

$$Y(c,e) = \begin{pmatrix} A(c,e) & 0 & \dots & \dots & 0 & 0 \\ 0 & A(c,e) & \dots & \dots & 0 & 0 \\ 0 & 0 & A(c,e) & \ddots & \vdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & 0 & A(c,e) \end{pmatrix},$$

where 0 denote the all zero  $n \times n$  matrix. Let  $\Omega = (\Omega_1, \dots, \Omega_m)$ , then (7) can be written as

$$P(c \rightarrow e | \alpha_i^j, i=1,2,\dots,n, j=1,2,\dots,m) \leq \exp(-\Omega Y(c,e) \Omega^* E_s / 4N_0). \quad (15)$$

Let  $\Theta$  denote the correlation of  $\Omega$ . Assume that  $\Theta$  is full rank (this is a physically acceptable assumption). The matrix  $\Theta$  being a non-negative definite square Hermitian matrix has a full rank  $nm \times nm$  lower triangular matrix  $C$  as it's square root. The diagonal elements of  $\Theta$  are unity, so that the rows of  $C$  are of length one. Define

$$v = (\varepsilon_1^1, \dots, \varepsilon_n^1, \varepsilon_1^2, \dots, \varepsilon_n^2, \dots, \varepsilon_1^m, \dots, \varepsilon_n^m)$$

by  $\Omega = vC^*$ , then it is easy to see that the components of  $v$  are uncorrelated complex Gaussian random variables with variance 0.5 per dimension. The mean of the components of  $v$  can be easily computed from the mean of  $a_i^j$  and the

26

matrix  $C$ . In particular of the  $a_i^j$  are of mean zero, so are the  $\varepsilon_i^j$ .

By (15), the conclusion is arrived at that

$$P(\mathbf{c} \rightarrow \mathbf{e} \mid \alpha_i^j, i=1,2,\dots,n, j=1,2,\dots,m) \leq \exp(-\gamma \mathbf{C}^* \mathbf{Y}(\mathbf{c}, \mathbf{e}) \mathbf{C} \gamma^* E_s / 4N_0) \quad (16)$$

5 The same argument can be followed as the case of independent fades with  $A(\mathbf{c}, \mathbf{e})$  replaced by  $\mathbf{C}^* \mathbf{Y}(\mathbf{c}, \mathbf{e}) \mathbf{C}$ . It follows that the rank of  $\mathbf{C}^* \mathbf{Y}(\mathbf{c}, \mathbf{e}) \mathbf{C}$  has to be maximized. Since  $C$  is full rank, this amounts to maximizing rank  $[Y(\mathbf{c}, \mathbf{e})] = m \text{ rank } [A(\mathbf{c}, \mathbf{e})]$ . Thus the rank criterion given  
10 for the independent fade coefficients holds in this case as well.

Since  $a_i^j$  are zero mean, so are  $\varepsilon_i^j$ . Thus by a similar argument to that of the case of independent fade coefficients, the conclusion that the determinant of  $\mathbf{C}^* \mathbf{Y}(\mathbf{c}, \mathbf{e}) \mathbf{C}$  must be maximized is arrived at. This is equal to  
15  $\det(\Theta) \det(Y(\mathbf{c}, \mathbf{e})) = \det(\Theta) [\det(A(\mathbf{c}, \mathbf{e}))]^m$ . In this light the determinant criterion given in the case of independent fade coefficients holds as well.

It follows from a similar argument that the rank  
20 criterion is also valid for the Rician case and that any code designed for Rician channel performs well for Rayleigh channel even if the fade coefficients are dependent. To obtain the gain criterion, one has to compute the mean of  $\varepsilon_i^j$  and apply the gain criterion given  
25 in the case of independent Rician fade coefficients. As appreciated by persons skilled in the art, this is a straightforward but tedious computation.

#### N. Space-Time Code Construction

30 In this section, the results of the previous section are used to design codes for a wireless communication system that employs  $n$  transmit antennas and (optional) receive antenna diversity, according to the second embodiment of the present invention.

27

The designed codes can be either trellis codes, or block codes having a trellis representation. Examples are provided of trellis codes, as generalization to block codes is straightforward, to persons skilled in the art.

5           i. Trellis Codes

In the design of the codes to be applied in the second illustrative embodiment, reference is made to those having the property that each transition branch at time  $t$  is labeled with a sequence  $q_t^1 q_t^2 \dots q_t^n$  of  $n$  symbols from the constellation alphabet  $Q$  for all  $1 \leq t \leq l$ . Any time that the encoder's path goes through such a transition branch, the symbol  $q_t^i$  is sent via antenna  $i$  for all  $1 \leq i \leq n$ .

The encoding for trellis codes is straightforward, with the exception that it is required that at the beginning and the end of each frame, the encoder be in known states. A method of decoding is illustrated next. Assuming channel estimates  $\hat{\alpha}_i^j$  of  $\alpha_i^j, i=1,2,\dots,n, j=1,2,\dots,m$  are available to the decoder. Assuming that  $r_t^i$  is the received signal at receive antenna  $i$  at time  $t$ , the decoder computes for any transition branch at time  $t$  having the label  $q_t^1 q_t^2 \dots q_t^n$ , the branch metric

$$\sum_{j=1}^m \left| r_t^j - \sum_{i=1}^n \hat{\alpha}_i^j q_t^i \right|^2.$$

The Viterbi algorithm is then used to compute the path with lowest accumulated metric.

25           The aforementioned trellis codes are called *Space-Time* codes, as they combine spatial and temporal diversity techniques. Furthermore, if the Space-Time code guarantee a diversity gain of  $rm$  for the multiple antenna communication systems discussed above, it is said that it is an  $r$ -Space-Time code.

30           A 4-state code for the 4-PSK constellation is given in Figur 9. For further illustration, there is also provided an 8-state code for the 8-PSK constellation

in Figure 11, and 8, 16, and 32-state codes for the 4-PSK constellation in Figures 12(a), 12(b), and 13, respectively. Also provided is a 16-state code for 16-QAM constellation in Figure 14.

5 The design rules that guarantee diversity for all the codes in Figures 11, 12(a), 12(b), 13, and 14 are:

- Transitions departing from the same state differ in the second symbol.
- Transitions arriving at the same state differ in the first symbol.

10

r-space-times for  $r \geq 2$ : As an illustration, a 4-space-time code for a 4 transmit antenna mobile communication systems is constructed. The input to the encoder is a block of length 2 of binary numbers corresponding to an integer  $i$  in  $Z_4 = \{0,1,2,3\}$ . The states of the trellis correspond to set of all three tuples  $(s_1, s_2, s_3)$  with  $s_i$  in  $Z_4$  for  $1 \leq i \leq 3$ . At state  $(s_1, s_2, s_3)$  upon receipt of input data  $i$ , the encoder outputs  $(i, s_1, s_2, s_3)$  elements of 4-PSK constellation (see Fig. 4(a)) and moves to state  $(i, s_1, s_2)$ . The performance of this code for 1 and 2 receive antennas is given in Figure 15.

15

20

#### 0. Channel Estimation and Interpolation

In both foregoing illustrated embodiments of the invention, it was assumed that the channel state information which is needed for decoding is known. However, in reality the receiver must estimate the channel state information. Also, the receiver must update this information as the channel varies. As illustrated in Figure 16, this may be accomplished by the periodic transmission of a probe or pilot symbol  $P$ , whose identity is known at the transmitting and the receiving sides of the communication apparatus.

30

During the transmission of the pilot symbols, the receiver derives estimates of the fade coefficients. The receiver estimates the channel state over the whole frame

35

of data using a channel interpolation scheme. The results of interpolation are used by the space-time decoder using decoding techniques known to the persons skilled in the art.

5           The inventors have observed that in high mobility environments inaccuracies in channel estimation and interpolation causes only a small number of errors in frames of data output by the space-time decoder. These few errors can be corrected using any outer block codes as  
10           are well-known in the art.

          Here is described an implementation for a wireless modem that employs the use of space-time codes according to the invention, along with a coding strategy called concatenated space-time coding.

15           P.       Basic Modem Architecture

          In this section the basic functions of a modem based on space-time coded modulation according to the invention are described. For the purpose of illustration, the channelization of the North American digital cellular  
20           standard IS-136 is assumed. However, the same modem architecture can be easily adopted to other channelization and/or any other application with minor modifications known to people skilled in the art.

          A brief overview of the frame structure in IS-136  
25           is as follows. On each 30 kHz wireless channel, the IS-136 standard defines 25 frames of data per second, each of which is then further subdivided into 6 time slots. Each time slot is of a 6.667 ms duration and carries 162 modulation symbols (modulation symbol rate is 24,300  
30           symbols/sec).

          i).   Transmitter

          Figure 17 shows a block diagram for a transmitter that employs space-time coding and is equipped with 2 transmit antennas (the extension of the same architecture  
35           to more than 2 transmit antennas is straightforward). A bit stream from the information source (either speech or data) is fed to the space-time encoder. The space-time

30

encoder groups each  $b$  information bits into one modulation symbol, where the number of bits  $b$  per modulation symbols will depend on the constellation used, which is assumed to be either M-QAM or M-PSK constellation. The space-time  
5 encoder uses a space-time code constructed according to criterion mentioned above.

Each group of  $b$  information bits generates two modulation symbols at the output of the space-time encoder. Each stream of modulation symbols is interleaved  
10 using a block interleaver. It is assumed that both bursts are interleaved in a similar way. Overhead, synchronization, and pilot symbols are then added to the output of each interleaver to build a burst. Each burst is then pulse-shaped using any suitable pulse shape known  
15 to persons skilled in the art, and transmitted from its corresponding antenna.

ii). Time Slot Structure

For the purpose of illustration, Figure 16 shows a slot structure for the case when the transmitter is  
20 equipped with two transmit antennas and follows IS-136 channelization. As mentioned, this slot structure can be easily extended to conform to other channelization and any number of transmit antennas.

In each time slot, two bursts are transmitted, one  
25 from each antenna. As in IS-136 North American Digital Cellular Standard, it is assumed that the modulation symbol rate is 24,300 symbols/sec and each burst consists of 162 symbols. Each burst starts with a 14 symbol synchronization sequence  $S_1$  and  $S_2$  that is used for timing  
30 and frequency synchronization at the receiver. In addition, the transmitter inserts 6 two-symbol pilot sequences  $P_1$  and  $P_2$  that will be used at the receiver to estimate the channel. The signal received at the receiver is the superposition of the two transmitted bursts, and in  
35 order to separate the two bursts at the receiver, it is necessary to define the two sequences  $S_1$  and  $S_2$  as well as the pilot sequences  $P_1$  and  $P_2$  as orthogonal sequences. It



is assumed that the synchronization and pilot symbols have the same energy per symbol as the information symbols. In addition, for the synchronization and pilot sequences  $\pi/4$ -shifted DQPSK modulation is used. Each burst will then have 136 symbols of information. The block interleaver will be then a  $17 \times 8$  block interleaver.

iii). Receiver

Figure 18 shows the corresponding block diagram for a mobile receiver equipped with two receive antennas according to this embodiment. For each receiver antenna, after matched filtering, the receiver splits the output samples into two streams.

The first stream contains the received samples that correspond to the information symbols. The second stream contains the received samples corresponding to the pilot symbols. These samples are first correlated with the pilot sequence for bursts transmitted from transmit antenna 1 to get an estimate for the channel (at the pilot positions) from transmit antenna 1 to the corresponding receive antenna. Also, the same set of samples are correlated with the pilot sequence for bursts transmitted from transmit antenna 2 to get an estimate for the channel (at the pilot positions) from transmit antenna 2 to the corresponding receive antenna. These estimates are then interpolated to form an estimate for channel state information needed for maximum likelihood decoding according to the metric previously defined. The interpolation filter can be designed in many ways known to persons skilled in the art. For optimum interpolation, a different interpolation filter should be used for every value of Doppler spread  $f_d$ , frequency offset  $f_o$ , and signal to noise ratio SNR. However this approach will be of great complexity for any practical implementation. Various approaches are proposed here. The first is to use a robust filter that will cover all possible range of operation, although this will lead to a slight degradation

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in performance at low Doppler and/or frequency offset values.

The second approach is to divide the range of operation into different regions, and for every region design an optimum interpolator for some operating point in that region and use that filter for the whole. By observing the channel correlation from the channel estimates or by observing the symbol error rate, the receiver can decide which filter to use for interpolation.

In addition, in estimating the channel over any burst, the pilot symbols in that burst are only used. This will minimize the overall system delay by avoiding the need to wait for future bursts in order to estimate the channel. Both data streams are then deinterleaved and fed to a vector Viterbi decoder.

iv). Basic Modem Performance

In this section, simulation results for the basic modem and time slot structure described above are presented. In addition, the pulse shape that was used is a square-root raised-cosine Nyquist pulse with roll-off factor of 0.35. At the receiver an oversampling factor of 8 is assumed.

Figure 19 shows the frame error rate (FER)  $P_f$  performance of the above modem for different values of Doppler spread  $f_d$  assuming perfect timing and frequency synchronization. For the static case, perfect knowledge of the CSI for comparison, is assumed. Plotted is  $P_f$  versus SNR (or symbol energy to noise ratio)  $E_s/N_0$ . For the ideal CSI, it can be seen that a FER of 0.1 at  $E_s/N_0$  of 14.75 dB. However, for a Doppler spread  $f_d$  of 170 Hz, which corresponds to a vehicle speed of 60 mph, the 0.1 FER is achieved at 20.5 dB  $E_s/N_0$ . For  $f_d = 120$  Hz, this number drops to 17.1 dB. It can also be noticed that FER floors at high  $E_s/N_0$ . In general, this increase in the required  $E_s/N_0$  as compared to the case with ideal CSI and the FER flooring are due to the errors in channel estimation.

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Q. Third Illustrative Embodiment-  
Concatenated Space-Time Coding

Figure 20 shows the distribution of the number of symbol errors per frame for the  $f_d = 170$  Hz for different values of  $E_s/N_0$ . For relatively high values of  $E_s/N_0$  ( $> 15$  dB), approximately 90% of the frames that are in error, the error is due to 5 symbol errors or less. Most of these errors can be recovered from, by concatenating the space-time code with any block code known to the persons skilled in the art, such as a Reed Solomon (RS) code. This overall coding strategy is designated concatenated space-time coding and is shown in Figures 21 and 22. Depending on the desired error correction capability and rate of the code and the type of constellation used, the dimension of the block code used should be such to produce an integer multiple of modulation symbols for each RS symbol. In this way, it will be possible to decode a burst immediately without the need to wait for other bursts and, thereby, minimize decoding delay. In addition, in this way, any symbol error at the output of the ST decoder will affect only one RS code symbol.

R. Modem Performance with  
Concatenated Space-Time Coding

The inventors simulated the above-described modem with the space-time code concatenated with a Reed-Solomon code. Three different shortened RS codes over  $GF(2^8)$  were used in the simulation. The first code, referred to as RS, is a shortened (68, 66) code that corrects single byte errors. The 66  $GF(2^8)$  symbols are first created by partitioning the bit stream into 66 groups of 8 bits each. The output 68  $GF(2^8)$  symbols are then partitioned into 136 16-QAM symbols, 2 channel symbols per one Reed-Solomon symbol, which are then fed to the ST encoder. The second code, referred to as RS3, is a shortened (68, 62) code that corrects three byte errors, and the third code, referred to as RS5, is a shortened (68, 58) code that corrects 5 byte errors. In this simulation, a timing

error of  $\pm T/16$  and a frequency offset  $f_o$  of 200 Hz are assumed.

Figure 23 shows the FER performance with concatenated space-time coding and in the presence of timing error and frequency offset for  $f_d = 170$  Hz. From this figure, it can be seen that in the case of the ST code alone a  $E_s/N_0$  of 23 dB is required to achieve  $P_F$  of 0.1. However, when the ST code is concatenated with RS3, for example, the required  $E_s/N_0$  is 17.5 dB, i.e., a 5.5 dB gain over the ST code alone. However, in this case, the net bit rate (per 30 kHz channel) will be reduced from 81.6 kbits/sec to 74.4 kbits/sec. If RS5 is used, the required  $E_s/N_0$  for  $P_F = 0.1$  will drop to 16.5 dB, which is only 1.75 dB higher than the case when ideal CSI are available. In this case, the net bit rate will be 69.6 kbits/sec.

#### S. Fourth Illustrative Embodiment- Multi Level Structured Space Time Codes

Some of the space-time codes described in the second embodiment of this invention may have multilevel structure. On occasions, it may be desirable to take advantage of this structure in practical communication systems, particularly when the number of transmit antennas is high. This has the significant advantage of reducing decoding complexity. Multilevel code structures and associated decoding techniques are known in the art. They can be combined with space-time coding, giving rise to the invention of a novel technique called multi-level space-time coding.

Without loss of generality, assume that the signal constellation  $Q_0$  consists of  $2^{b_0}$  signal points. Assume that  $f$ -levels of coding is used. Associated with this  $f$ -levels of coding, a partition based on subsets

$$Q_{f-1} \subset Q_{f-2} \subset \dots \subset Q_1 \subset Q_0$$

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is chosen with the number of elements of  $Q_j$  equal to  $2^{b_j}$  for all  $0 \leq j \leq f-1$ . By such a partitioning, it is meant that  $Q_0$  is the union  $2^{b_0-b_1}$  disjoint sets called cosets of  $Q_1$  in  $Q_0$ , each having  $2^{b_1}$  elements such that one of these

5 cosets is  $Q_1$ . Having the cosets of  $Q_1$  in  $Q_0$  at hand, each coset is then divided into  $2^{b_1-b_2}$  disjoint sets each having  $2^{b_2}$  elements. The  $2^{b_1-b_2}$  subsets of  $Q_1$  are called cosets of  $Q_2$  in  $Q_1$ . The set of cosets of  $Q_2$  in  $Q_1$  must include  $Q_2$ . Thus there are  $2^{b_0-b_2}$  subsets of  $Q_0$  with  $2^{b_2}$  elements called

10 the cosets of  $Q_2$  in  $Q_0$ . The set of cosets of  $Q_2$  in  $Q_0$  includes  $Q_2$ . This procedure is repeated until cosets of  $Q_j$  in  $Q_k$  for all  $0 \leq k < j \leq f-1$  are arrived at. Let  $r_{f-1}=b_{f-1}$  and  $r_j=b_{j+1}-b_j$  for  $j=0,1,\dots,f-2$ . Then  $Q_j$  contains  $2^{r_j}$  cosets of  $Q_{j+1}$  for all  $j=0,1,\dots,f-2$ .

15 Every  $K=K_0+\dots+K_{f-1}$  bits of input data is encoded using encoders  $0,1,\dots,f-1$  corresponding to the  $f$  levels. It is required that all the encoders have a trellis representation. At each time  $t$  depending on the state of the  $j$ -th encoder and the input data, a branch of the

20 trellis of the  $j$ -th encoder is chosen which is labeled with  $n$  blocks of  $r_j$  bits denoted by  $B_t^i(j), B_t^i(j), \dots, B_t^i(j)$ . The blocks  $B_t^i(0), \dots, B_t^i(f-1)$  then choose a point of the signal constellation in the following way.

The block  $B_t^i(0)$  chooses a coset  $Q_1^i$  of  $Q_1$  in  $Q_0$ .

25 The block  $B_t^i(1)$  chooses a coset  $Q_2^i$  of  $Q_2$  in  $Q_1$  and so forth. Finally the block  $B_t^i(f-1)$  chooses a point of  $Q_{f-1}^i$  a coset of  $Q_{f-1}$  chosen in the last step. The chosen point is then transmitted using the  $i$ -th antenna for  $1 \leq i \leq n$ . Multilevel decoding can be done in a manner known to

30 those skilled in the art.

Suppose that the encoder of the  $j$ -th level has  $2^{s_j}$  states at time  $t$ . One can view the multi-level code

described above as a space-time code  $C$  with  $2^{(s_0 + \dots + s_{f-1})}$  states at time  $t$ . The states of  $C$  correspond  $f$ -tuples  $(s_t^0, s_t^1, \dots, s_t^{f-1})$  of states of encoders  $0, 1, \dots, f-1$ . The branch labels between states  $(s_t^0, s_t^1, \dots, s_t^{f-1})$  and  $(s_{t+1}^0, s_{t+1}^1, \dots, s_{t+1}^{f-1})$  is the set of symbols that are sent via antennas  $1, 2, \dots, n$  if each encoder  $j$  goes from states  $s_t^j$  to the state  $s_{t+1}^j$  for  $0 \leq j \leq f-1$ . In this way, one can view a multi-level space-time code as a regular space-time code with a multi-level structure that allows simplified decoding. The penalty for this simplified decoding is a loss in performance. Also, the design criterion derived previously could be applied to the space-time code  $C$ . Alternatively the design criteria can instead be applied to the trellis of each encoder  $0 \leq j \leq f-1$  providing different diversities at each level.

The discussion of the illustrative embodiment above is illustrated with an example. Consider the transmission of 4-bits/sec/HZ using the 16-QAM constellation and the set partitioning of Figure 24. At each time input bits are grouped into two blocks of two bits. The first and second blocks of two bits input data are respectively the input to the first and second encoder whose trellis is given in Figure 25. Each branch of this trellis is labeled with two blocks of two bits of data. These two bits are represented with numbers 0, 1, 2 and 3. Upon the choice of branches with respective labels  $a_1 a_2$  and  $b_1 b_2$  by the zero-th and the first encoders, the signal points  $4a_1 + b_1$  and  $4a_2 + b_2$  are sent via antennas 1 and 2. The equivalent 16-state space-time trellis code is given in Figure 26.

#### T. Fifth Illustrative Embodiment: Smart-Greedy Codes

Smart greedy codes are a class of space-time codes of particular interest in the implementation of the invention. These codes are able to take special advantage of possible rapid changes in the channel

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without any feedback from the receiver. The idea is to construct codes using a hybrid criteria such that possible rapid changes in the channel is taken into account by the design criteria. In this light, an analysis is provided for the case of rapidly fading channels as well.

i) Analysis of Rapid Fading

In this connection, the model of a mobile communication system having  $n$  antennas at the base and  $m$  antennas at the mobile station is refined. Data is encoded using a channel code. As in other embodiments, the encoded data goes through a serial to parallel device and is divided into  $n$  streams of data. Each stream of data is used as the input to a pulse shaper. The output of each shaper is then modulated using a modulator. At each time the output of modulator  $i$  is a signal that is transmitted using transmit antenna (Tx antenna)  $i$  for  $1 \leq i \leq n$ . Again, the  $n$  signals are transmitted simultaneously each from a different transmit antenna and all these signals have the same transmission period  $T$ . The signal at each receive antenna is a noisy version of the superposition of the faded versions of the  $n$  transmitted signals. Assume that each element of the signal constellation is contracted by a scale factor  $\sqrt{E_s}$ , chosen so that the average energy of the constellation elements is 1.

At the receiver, the demodulator makes decision statistic based on the received signals at each receive antenna  $1 \leq j \leq m$ . Let  $c_t^i$  denote the transmitted symbol from the  $i$ -th transmit antenna at transmission interval  $t$  and  $d_t^j$  be the receive word at the receive antenna  $j$ . Then,

$$d_t^j = \sum_{i=1}^n \alpha_i^j(t) c_t^i \sqrt{E_s} + \eta_t^j \quad (18)$$

This is equivalent to the <sup>38</sup> assumption that signals transmitted from different antennas undergo independent fades. The coefficients  $\alpha_i^j(t)$  are modeled as samples of a stationary complex Gaussian stochastic process with mean zero and variance 0.5 per dimension. Also,  $\eta_k^j$  are independent samples of a zero mean complex white Gaussian process with two sided power spectral density  $N_0/2$  per dimension. For the static fading case, suppose that  $\alpha_i^j(t)$  are constant during a frame and are independent from one frame to another and a design criterion was established. When the fading is rapid, the coefficients  $\alpha_i^j(t), t=1,2,\dots,l, i=1,2,\dots,n, j=1,2,\dots,m$  are modeled as independent samples of a complex Gaussian process with mean zero and variance 0.5 per dimension, and another design criteria is established as follows.

Assuming that the coefficients  $\alpha_i^j(t)$  for  $t=1,2,\dots,l, i=1,2,\dots,n, j=1,2,\dots,m$  are known to the decoder, the probability of transmitting

$$\mathbf{c} = c_1^1 c_1^2 \dots c_1^n c_2^1 c_2^2 \dots c_2^n \dots c_l^1 c_l^2 \dots c_l^n,$$

and deciding in favor of

$$\mathbf{e} = e_1^1 e_1^2 \dots e_1^n e_2^1 e_2^2 \dots e_2^n \dots e_l^1 e_l^2 \dots e_l^n$$

at the decoder is well approximated by

$$P(\mathbf{c} \rightarrow \mathbf{e} \mid \alpha_i^j, i=1,2,\dots,n, j=1,2,\dots,m, t=1,2,\dots,l) \leq \exp(-d^2(\mathbf{c}, \mathbf{e}) E_s / 4N_0)$$

where

$$d^2(\mathbf{c}, \mathbf{e}) = \sum_{t=1}^l \sum_{i=1}^n \left| \sum_{j=1}^m \alpha_i^j(t) (c_i^j - e_i^j) \right|^2 \quad (19)$$

This is the standard approximation to the Gaussian tail function.

Let

$$\Omega_i(t) = (\alpha_i^1(t), \alpha_i^2(t), \dots, \alpha_i^m(t))$$



and  $C(t)$  denote the  $n \times n$  matrix with the element at  $p$ -th row and  $q$ -th column equal to  $(c_t^p - e_t^p)(\bar{c}_t^q - \bar{e}_t^q)$ . Then it can be seen that

$$d^2(c, e) = \sum_{p=1}^n \sum_{q=1}^n \Omega_p(t) C(t) \Omega_q^*(t) \quad (20)$$

5 The matrix  $C(t)$  is Hermitian, thus there exist a unitary matrix  $V(t)$  and a diagonal matrix  $D(t)$  such that  $C(t) = V(t)D(t)V^*(t)$ . The diagonal elements of  $D(t)$ , denoted here by  $D_{ii}(t)$ ,  $1 \leq i \leq n$ , are the eigenvalues of  $C(t)$  counting multiplicities. Since  $C(t)$  is Hermitian, these  
10 eigenvalues are real numbers. Let

$$\Lambda_j(t) = \Omega_j(t)V(t) = (\lambda_j^1(t), \dots, \lambda_j^n(t)),$$

then  $\lambda_j^i(t)$  for  $i=1, 2, \dots, n, j=1, 2, \dots, m, t=1, 2, \dots, l$  are independent complex Gaussian variables with mean zero and variance 0.5 per dimension and

$$15 \quad \Omega_j(t) C(t) \Omega_j^*(t) = \sum_{i=1}^n |\lambda_j^i(t)|^2 D_{ii}(t).$$

By combining this with (19) and (20) and averaging with respect to the Rayleigh distribution of  $|\lambda_j^i(t)|$ , the following is arrived at

$$P(c \rightarrow e) \leq \prod_{i,t} (1 + D_{ii}(t) \frac{E_s}{4N_0})^{-m}. \quad (21)$$

20 The matrix  $C(t)$  is next examined. The columns of  $C(t)$  are all different multiples of

$$c_t - e_t = (c_t^1 - e_t^1, c_t^2 - e_t^2, \dots, c_t^n - e_t^n).$$

Thus,  $C(t)$  has rank 1 if  $c_t^1 c_t^2 \dots c_t^n \neq e_t^1 e_t^2 \dots e_t^n$  and rank zero otherwise. It follows that  $n-1$  elements in the list

$$25 \quad D_{11}(t), D_{22}(t), \dots, D_{nn}(t)$$

are zeros and the only possible nonzero element in this list is  $|c_t - e_t|^2$ . By (21), it can now be concluded that

$$P(c \rightarrow e) \leq \prod_{t=1}^l \left( 1 + |c_t - e_t| \frac{E_s}{4N_0} \right)^{-m} \quad (22)$$

40

Let  $V(c, e)$  denote the set of time instances  $1 \leq t \leq l$  such that  $|c_t - e_t| \neq 0$  and let  $|V(c, e)|$  denote the number of elements of  $v(c, e)$ . Then it follows from (22) that

$$P(c \rightarrow e) \leq \prod_{t \in V(c, e)} \left( |c_t - e_t| \frac{E_s}{4N_0} \right)^{-m}. \quad (23)$$

- 5 It follows that a diversity of  $m|V(c, e)|$  is achieved. Examining the coefficient of  $(E_s/4N_0)^{-m|V(c, e)|}$  leads to the desired design criterion. Below, this criterion is combined with that of static flat fading case given before to arrive at a hybrid criteria.

10 U. A Hybrid Design Criteria For  
For Smart Greedy Space-Time  
Codes:

The Distance/Rank Criterion: In order to achieve the diversity  $\alpha m$  in a rapid fading environment, for any  
15 two codewords  $c$  and  $e$  the strings  $c_1^1 c_1^2 \dots c_1^n$  and  $e_1^1 e_1^2 \dots e_1^n$  must be different at least for  $\alpha$  values of  $1 \leq t \leq n$ . Furthermore, let

$$B(c, e) = \begin{pmatrix} e_1^1 - c_1^1 & e_1^2 - c_1^2 & \dots & \dots & e_1^n - c_1^n \\ e_2^1 - c_2^1 & e_2^2 - c_2^2 & \dots & \dots & e_2^n - c_2^n \\ e_3^1 - c_3^1 & e_3^2 - c_3^2 & \ddots & & e_3^n - c_3^n \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ e_n^1 - c_n^1 & e_n^2 - c_n^2 & \dots & \dots & e_n^n - c_n^n \end{pmatrix}$$

If  $B(c, e)$  has minimum rank  $r$  over the set of pairs of  
20 distinct codeword, then a diversity of  $\alpha m$  is achieved in static flat fading environments.

The Product/Determinant Criterion: Let  $V(c, e)$  denote the set of time instances  $1 \leq t \leq l$  such that  $c_1^1 c_1^2 \dots c_1^n \neq e_1^1 e_1^2 \dots e_1^n$  and let  $|c_t - e_t|^2 = \sum_{i=1}^n |c_t^i - e_t^i|^2$ . Then to  
25 achieve the most coding gain in a rapid fading environment, the minimum of the products  $\prod_{t \in V(c, e)} |c_t - e_t|^2$  taken over distinct codewords  $e$  and  $c$  must be maximized. For the case of a static fading channel, the minimum of  $r$ -th roots of the sum of determinants of all  $r \times r$

principal cofactors of  $A(c, e) = B(c, e)B^*(, e)$  taken over all pairs of distinct codewords  $e$  and  $c$  corresponds to the coding gain, where  $r$  is the rank of  $A(c, e)$ .

5 The construction of illustrative implementations of smart greedy codes according to this embodiment of the invention is illustrated with some examples. It will be assumed that at the beginning and the end of the frame, the encoder is in the zero state.

Example A: Suppose that a transmission rate of 10 0.5 bits/sec/Hz is required. In this example and as illustrated in Fig. 27(a), the BPSK constellation is used, with 0 denoting  $\sqrt{E_s}$  and 1 denoting  $-\sqrt{E_s}$ . The objective is to guarantee diversity gains 2 and 4 respectively in slow and rapid flat fading environments. 15 The following code using M-TCM construction guarantees these diversity gains. At any time  $2k + 1, k = 0, 1, 2, \dots$  depending on the state of the encoder and the input bit a branch is chosen by the encoder and the first coordinate and second coordinates of the labels are sent 20 simultaneously from Tx antennas at times  $2k + 1$  and  $2k + 2$ . For instance at time 1, if the branch label 10 11 is chosen, symbols 1,0 and 1,1 are sent respectively from transmit antennas one and two at times one and two.

Example B: Here a transmission rate of 1 25 bits/sec/Hz and diversity gains of 2 and 3 respectively in static and rapid flat fading environments are desired. In this example, illustrated in Fig. 27(b), the 4-PSK constellation is used instead. The objective is to guarantee diversity gains 2 and 3 respectively in slow 30 and rapid flat fading environments. The following code using M-TCM construction guarantees these diversity gains. At times  $t = 3k, k = 0, 1, 2, \dots$ , three bits of data arrive at the encoder. The first bit choose a branch depending on the state of the encoder and the rest of two 35 bits choose one of the 4 labels of that branch such as  $b_t^1 b_t^2 b_{t+1}^1 b_{t+1}^2 b_{t+2}^1 b_{t+2}^2$ . Then  $b_t^1, b_{t+1}^1$  and  $b_{t+2}^1$  are sent via antenna

1 respectively at times  $t, t + 1$  and  $t + 2$ . Similarly,  
 $b_t^2, b_{t+1}^2$  and  $b_{t+2}^2$  are sent via antenna 2 respectively at  
time  $t, t + 1$  and  $t + 2$ .

5 As before, the inventors have simulated the  
performance of communication systems designed based on  
the above code. Excellent results have been confirmed in  
both fast and slow fading environments.

\* \* \*

10

The foregoing description of the system and  
method of the invention is illustrative, and variations  
in construction and implementation will occur to persons  
skilled in the art. For example, although the present  
15 invention is described in the time domain, frequency  
domain analogs or variants of it easily occur to those  
skilled in the art. For instance, space-time codes  
presented in the second illustrative embodiment can be  
easily applied to DS-CDMA communication systems. To  
20 illustrate, assume that user X is provided with two  
transmit antennas (with generalization to  $n$  antennas  
being trivial to those skilled in the art). User X  
chooses a space-time code designed to be used with two  
transmit antennas. User X can use a similar PN sequence  
25 for data transmission from both antennas. At the  
receiver side, correlation with the aforementioned  
sequence gives a sum of faded versions of the signals  
transmitted from each antenna. In this light, decoding  
of the space-time code can be carried out in a manner  
30 similar to those described in the second embodiment of  
this work as well.

Alternatively, user X can use distinct PN  
sequences for transmission from both transmit antennas.  
If the PN sequences used to transmit from both antennas  
35 are orthogonal to each other, at the receiver co-relation  
with the first or second sequence gives respectively

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noisy versions of the transmitted signals from antennas one or two, which can be used for decoding at the receiver. This has a penalty in terms of bandwidth expansion but can be used to increase the data rate and/or provide diversity advantage.

In general, it is also possible to choose two arbitrary PN sequences for two transmit antennas. Correlation with these sequences at the receiver side gives sums of faded versions of multiples of the transmitted signals that can be used for decoding.

The above discussion demonstrates a DS-CDMA analog of the space-time coding. Analogs of the embodiments of the present invention in frequency domain also can easily be obtained, but are not discussed here.

For further instance, while mobile cellular implementations have been described, the invention could be applied to other communication environments. The invention is accordingly intended to be limited only by the following claims.

WHAT IS CLAIMED IS:

1. A method of transmitting a digital signal using a plurality of  $n$  antennas, comprising the steps of:  
 selecting a code, the code to be used by an  
 5 encoder to generate the digital signal, and having a trellis structure such that the branches of the trellis are elements of a signal constellation;  
 selecting a branch of the trellis based on input data, and a previous state of the encoder; and  
 10 transmitting elements of the signal constellation corresponding to a branch label of the selected branch over the plurality of antennas simultaneously.
2. The method of claim 1, wherein the step of  
 15 selecting a code comprises the step of selecting a code having a trellis in which:  
 the label of branches departing a state differ in at least one location; and  
 the label of branches arriving in that state  
 20 differ in at least one location.
3. The method of claim 2, wherein:  
 the step of selecting a branch of the trellis comprises the step of selecting, at each time  $t$ , branches of the form  $q_t^1 q_t^2 \dots q_t^n$ ; and  
 25 the step of transmitting comprises the step of transmitting, at each time  $t$ , the signals  $q_t^i, i=1,2,\dots,n$  over antennas  $1 \leq i \leq n$  simultaneously.
4. The method of claim 1 wherein the trellis code has the properties:  
 30 given any two distinct paths  $P_1$  and  $P_2$  emerging from a state of the trellis at time  $t_1$  and re-emerging in another state of the trellis at time  $t_2 + 1$ , where  $P_1$  and  $P_2$  correspond respectively to branch labels  
 $c_{t_1}^1 c_{t_1}^2 \dots c_{t_1}^n, \dots, c_{t_2}^1 c_{t_2}^2 \dots c_{t_2}^n$  and  $e_{t_1}^1 e_{t_1}^2 \dots e_{t_1}^n, \dots$ , and  $e_{t_2}^1 e_{t_2}^2 \dots e_{t_2}^n$ , the  
 35 matrix

$$B(P_1, P_2) = \begin{pmatrix} e_{t_1}^1 - c_{t_1}^1 & \dots & e_{t_2}^1 - c_{t_2}^1 \\ e_{t_1}^2 - c_{t_1}^2 & \dots & e_{t_2}^2 - c_{t_2}^2 \\ e_{t_1}^3 - c_{t_1}^3 & \ddots & e_{t_2}^3 - c_{t_2}^3 \\ \vdots & \ddots & \vdots \\ e_{t_1}^n - c_{t_1}^n & \dots & e_{t_2}^n - c_{t_2}^n \end{pmatrix}$$

has rank greater or equal to  $r$  for some positive integer  $r$ , and

5 if  $r = n$ , the determinant of  $B(P_1, P_2)B^*(P_1, P_2)$  is greater than some non-negative number  $b$ , where  $B^*(P_1, P_2)$  is the conjugate transpose of  $B(P_1, P_2)$ .

5. The method of claim 4, wherein  $r$  is greater than a predetermined amount.

10 6. The method of claim 4, wherein  $r = n$  and  $b$  is made greater than a predetermined amount.

7. The method of claim 1, wherein the step of transmitting comprises the step of using quadrature and in-phase amplitude values based on a symbol.

15 8. The method of claim 1, wherein the step of transmitting comprises the step of using phase shifts.

9. The method of claim 1, further comprising the step of delaying the transmission of at least one symbol by a predetermined time delay.

20 10. The method of claim 1, wherein at least one of the branch labels is a zero signal.

11. The method of claim 1, further comprising the step of receiving the transmitted symbols on a plurality of  $m$  receive antennas.

25 12. The method of claim 1, wherein the step of selecting a code having a trellis structure further comprises the step of selecting a code having a trellis in which:

30 the minimum of the rank of  $B(P_1, P_2)$  over distinct paths  $P_1$  and  $P_2$  emerging from a state of the trellis at time  $t_1$  and remerging in another state of the trellis at time  $t_2 + 1$  is  $r$ , in which case, given a

plurality of  $m$  receive antennas at the receiver, a diversity of  $rm$  can be achieved.

13. The method of claim 1, wherein the step of selecting a code having a trellis structure further comprises the step of selecting a code having a trellis in which:

the minimum of the rank of  $B(P_1, P_2)$  over distinct paths  $P_1$  and  $P_2$  emerging from a state of a trellis at time  $t_1$  and remerging in another state of the trellis at time  $t_2 + 1$  is  $r$ , and the minimum of the sum of  $r \times r$  determinants of the principal cofactors of  $B(P_1, P_2)B^*(P_1, P_2)$  is greater than some non-negative number  $p$ .

14. The method of claim 13, wherein  $p$  is greater than a predetermined amount.

15. The method of claim 1, wherein the selected code is a space-time code.

16. A method of transmitting a digital signal using a plurality of  $n$  antennas, comprising the steps of:  
selecting a multi-level structured space-time code, with encoders at each level having a trellis representation;

selecting branches of the trellises at different levels, and selecting constellation points based on the label of the selected branches and some set partitioning of the constellation; and

transmitting the signal constellation points corresponding to the branch labels of the selected branches over the plurality of antennas simultaneously.

17. The method of claim 16, wherein the step of selecting a multi-leveled structured space-time code comprises the step of selecting a code having a trellis in which:

the label of branches departing a state differ in at least one location; and

the label of branches arriving in that state differ in at least one location.



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18. The method of claim 17, wherein:

the step of selecting branches of the trellis comprises the step of selecting, at each time  $t$ , branches of the form  $q_t^1, q_t^1, \dots, q_t^n$ ; and

5 the step of transmitting comprises the step of transmitting, at each time  $t$ , the signals  $q_t^i, i=1, 2, \dots, n$  over antennas  $1 \leq i \leq n$  simultaneously.

19. The method of claim 16, wherein the step of transmitting comprises the step of using quadrature and  
10 in-phase amplitude values based on a symbol.

20. The method of claim 16, wherein the step of transmitting comprises the step of using quadrature and in-phase amplitude values based on a symbol.

21. The method of claim 16, wherein the step of  
15 transmitting comprises using phase shifts.

22. The method of claim 16, further comprising the step of delaying the transmission of at least one symbol by a predetermined time delay.

23. The method of claim 16, wherein at least one  
20 of the transmitted signals is a zero signal.

24. The method of claim 16, further comprising a step of receiving the transmitted symbols on a plurality of  $m$  receive antennas.

25. The method of claim 16, wherein different  
25 diversities at different levels of the code are achieved.

26. The method of claim 16, wherein multilevel decoding is employed.

27. A method of transmitting a digital signal using a plurality of  $n$  antennas, comprising the steps of:

30 selecting a code, the code to be used by an encoder to generate the digital signal, and having a trellis structure, such that the branches of the trellis at time  $t = 1, t_0 + 1, 2t_0 + 1, \dots$  are of the form  $q_t^1 q_t^2 \dots q_t^n, q_{t+1}^1 q_{t+1}^2 \dots q_{t+1}^n, \dots, q_{t+t_0-1}^1 q_{t+t_0-1}^2 \dots q_{t+t_0-1}^n$  with  $q_t^i, i=1, 2, \dots, n, t=1, 2, \dots$   
35 elements of a signal constellation;

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selecting, at each time  $t$ , a branch of the trellis based on input data, and a previous state of the encoder;

dividing the label of the selected branch  
 5 into to blocks of length  $n$  of the signal constellation denoted by  $q_t^1 q_t^2 \dots q_t^n, q_{t+1}^1 q_{t+1}^2 \dots q_{t+1}^n, \dots, q_{t+t_0-1}^1 q_{t+t_0-1}^2 \dots q_{t+t_0-1}^n$ ; and

transmitting the signals  $q_{t+q}^i, i=1,2,\dots,n$ ,  
 using transmit antennas  $1 \leq i \leq n$  simultaneously at time  
 $t + q$  for  $0 \leq q \leq t_0-1$ .

10 28. The method of claim 27, wherein the step of selecting a code comprises the step of selecting a code having a trellis in which:

the label of branches departing a state differ in at least one location; and

15 the label of branches arriving in that state differ in at least one location.

29. The method of claim 27 wherein:

given any two distinct paths  $P_1$  and  $P_2$   
 emerging from a state of the trellis at a time and  
 20 remerging in another state of the trellis at another time, where  $P_1$  and  $P_2$  correspond respectively to branch labels  $c_{t_1}^1 c_{t_1}^2 \dots c_{t_1}^n, \dots, c_{t_2}^1 c_{t_2}^2 \dots c_{t_2}^n$  and  $e_{t_1}^1 e_{t_1}^2 \dots e_{t_1}^n, \dots$  and  $e_{t_2}^1 e_{t_2}^2 \dots e_{t_2}^n$ , the matrix

$$B(P_1, P_2) = \begin{pmatrix} e_{t_1}^1 - c_{t_1}^1 & \dots & e_{t_2}^1 - c_{t_2}^1 \\ e_{t_1}^2 - c_{t_1}^2 & \dots & e_{t_2}^2 - c_{t_2}^2 \\ e_{t_1}^3 - c_{t_1}^3 & \ddots & e_{t_2}^3 - c_{t_2}^3 \\ \vdots & \ddots & \vdots \\ e_{t_1}^n - c_{t_1}^n & \dots & e_{t_2}^n - c_{t_2}^n \end{pmatrix}$$

25 has rank greater or equal to  $r$  for some positive integer  $r$ , and the strings  $c_{t_1}^1 c_{t_1}^2 \dots c_{t_1}^n$  and  $e_{t_1}^1 e_{t_1}^2 \dots e_{t_1}^n$  are different for at least  $v$  values of  $t_1 \leq t \leq t_2$ ,

the minimum of  $r$ -th roots of the sum of  
 30 determinants of all  $r \times r$  principal cofactors of  $A(c, e) = B(c, e) B^*(c, e)$  taken over all pairs of distinct

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codewords  $\mathbf{e}$  and  $\mathbf{c}$  is greater than some non-negative number  $W$ ,

if  $r = n$ , the determinant of  $B(P_1, P_2)B^*(P_1, P_2)$  is greater than some non-negative number  $A$ , where  
 5  $B^*(P_1, P_2)$  is the conjugate transpose of  $B(P_1, P_2)$ , and

the product  $\prod_{t \in v(\mathbf{c}, \mathbf{e})} |\mathbf{c}_t - \mathbf{e}_t|^2$  taken over distinct codewords  $\mathbf{e}$  and  $\mathbf{c}$  is greater than a nonnegative number  $A'$ , where  $v(\mathbf{c}, \mathbf{e})$  is the set of time instances  $t_1 \leq t \leq t_2$  such that  $c_t^1 c_t^2 \cdots c_t^n \neq e_t^1 e_t^2 \cdots e_t^n$  and  
 10  $|\mathbf{c} - \mathbf{e}|^2 = \sum_{i=1}^n |c_i - e_i|^2$ .

30. The method of claim 29, wherein  $r$  is greater than a predetermined amount.

31. The method of claim 29, wherein  $v$  is greater than a predetermined amount.

15 32. The method of claim 29, wherein  $W$  is made greater than a predetermined amount.

33. The method of claim 29, wherein  $A'$  is greater than a predetermined amount.

20 34. The method of claim 29, wherein  $r = n$  and  $A$  is greater than a predetermined amount.

35. The method of claim 29, wherein the step of transmitting comprises the step of using quadrature and in-phase amplitude values based on a symbol.

25 36. The method of claim 29, wherein the step of transmitting comprises using phase shifts.

37. The method of claim 29, further comprising the step of delaying the transmission of at least one symbol by a predetermined time delay.

30 38. The method of claim 29, wherein at least one of the branch labels is a zero signal.

39. The method of claim 29, further comprising a step of receiving the transmitted symbols on a plurality of  $m$  receive antennas.

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40. The method of claim 29, wherein the selected code provides further diversity gain when the channel has rapid changes.

41. The method of claim 29, wherein the selected  
5 code is a smart greedy code.

42. A method of transmitting a digital signal using a plurality of n antennas, comprising the steps of:

encoding the digital signal using an outer error-correcting code;

10 encoding the output of the outer error correcting code using an inner space-time code;

transmitting the code symbols of the space-time code simultaneously using n transmit antennas; and

transmitting pilot signals during certain  
15 time intervals to test channel states.

43. The method of claim 42, wherein the outer code is one of a BCH code or a Reed-Solomon code.

44. The method of claim 42, wherein the space-time code has a trellis structure.

20 45. The method of claim 42, wherein the pilot symbol sequence transmitted from a transmit antenna is orthogonal to those of other antennas.

46. The method of claim 42, further comprising the steps of:

25 receiving the transmitted signals; and  
correlating the received signals at predetermined time intervals with the transmitted pilot sequences.

30 47. The method of claim 42, further comprising the step of interleaving the encoded data using at least one interleaver.

48. The method of claim 42, further comprising the step of interpolating channel states using at least one interpolator.

35 49. The method of claim 47, further comprising the step of permuting the transmitted symbols from different antennas in different manners.

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50. The method of claim 42, wherein the step of transmitting comprises the step of quadrature and in-phase amplitude values based on a symbol.

5 51. The method of claim 42, wherein the step of transmitting branch labels comprises using phase shifts.

52. The method of claim 42, comprising the step of delaying the transmission of at least one symbol by a predetermined time delay.

10 53. The method of claim 42, where at least one of the branch labels is a zero signal.

54. The method of claim 42, further comprising the step of receiving the transmitted symbols on a plurality of m receive antennas.

15 55. The method of claim 48, further comprising the step of deinterleaving using a plurality of deinterleavers.

56. The method of claim 42, wherein space-time decoding is followed by outer-code decoding.

20 57. An apparatus for transmitting a digital signal using a plurality of n antennas, comprising:

an encoder unit for selecting a code, the code to be used to generate the digital signal, and having a trellis structure such that the branches of the trellis are elements of a signal constellation, the  
25 encoder unit selecting a branch of the trellis based on input data, and a previous state of the encoder unit; and  
a transmitter unit, for transmitting elements of the signal constellation corresponding to a branch label of the selected branch over the plurality of  
30 antennas simultaneously.

58. The apparatus of claim 57, wherein the encoder unit selects a code having a trellis in which:

the label of branches departing a state differ in at least one location; and  
35 the label of branches arriving in that state differ in at least one location.

59. An apparatus for transmitting a digital signal using a plurality of n antennas, comprising:

a multi-level structured space-time code encoder unit, the encoder unit at each level having a trellis representation, and selecting:

branches of the trellises at different levels, and

constellation points based on labels of the selected branches; and

a transmitter unit, for transmitting the signal constellation points corresponding to the branch labels of the selected branches over the plurality of antennas simultaneously.

60. The apparatus of claim 59, wherein the multi-leveled structured space-time encoder unit selects a code having a trellis in which:

the label of branches departing a state differ in at least one location; and

the label of branches arriving in that state differ in at least one location.

61. An apparatus for transmitting a digital signal using a plurality of n antennas, comprising:

an encoder unit, the encoder unit selecting a code to be used by the encoder to generate the digital signal, the code having a trellis structure, such that the branches of the trellis at time  $t = 1, t_0 + 1, 2t_0 + 1, \dots$  are of the form

$$q_t^1 q_t^2 \dots q_t^n, q_{t_0+1}^1, q_{t_0+1}^2 \dots q_{t_0+1}^n, \dots, q_{t_0+t_0-1}^1 \dots q_{t_0+t_0-1}^n \text{ with}$$

$q_t^i, i=1, 2, \dots, n, t=1, 2, \dots$  elements of a signal constellation,

and the encoder unit:

selecting, at each time t, a branch of the trellis based on input data, and a previous state of the encoder unit, and

dividing the label of the selected branch  
 into to blocks of length  $n$  of the signal constellation  
 denoted by  $q_t^1 q_t^2 \dots q_t^n, q_{t+1}^1 q_{t+1}^2 \dots q_{t+1}^n, \dots, q_{t+t_0-1}^1 q_{t+t_0-1}^2 \dots q_{t+t_0-1}^n$ ; and

a transmitter unit, for transmitting the  
 5 signals  $q_{t,q}^i, i=1,2,\dots,n$ , using transmit antennas  $1 \leq i \leq n$   
 simultaneously at time  $t + q$  for  $0 \leq q \leq t_0-1$ .

62. The apparatus of claim 61, wherein the  
 encoder unit selects a code having a trellis in which:

the label of branches departing a state  
 10 differ in at least one location; and

the label of branches arriving in that state  
 differ in at least one location.

63. An apparatus for transmitting a digital  
 signal using a plurality of  $n$  antennas, comprising:

15 an encoder unit for encoding the digital  
 signal using an outer error-correcting code, and encoding  
 the output of the outer error correcting code using an  
 inner space-time code;

a transmitter unit, for transmitting:  
 20 the code symbols of the space-time code  
 simultaneously using  $n$  transmit antennas, and  
 pilot signals during certain time intervals  
 to test channel states.

64. The apparatus of claim 63, wherein the outer  
 25 code is one of a BCH code or a Reed-Solomon code.

65. The apparatus of claim 63, wherein the  
 space-time code has a trellis structure.

66. The apparatus of claim 63, wherein the pilot  
 symbol sequence transmitted from a transmit antenna is  
 30 orthogonal to those transmitted from other antennas.

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FIG. 1A

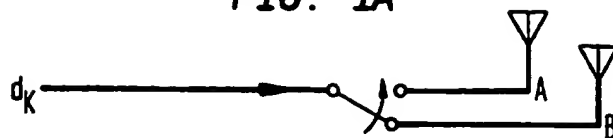


FIG. 1B

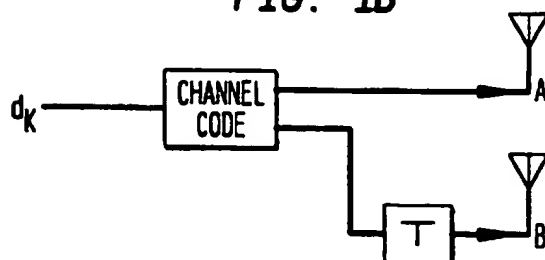


FIG. 1C

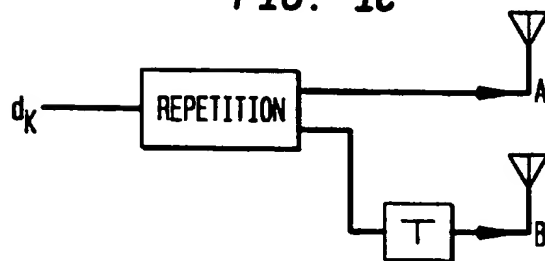


FIG. 2A

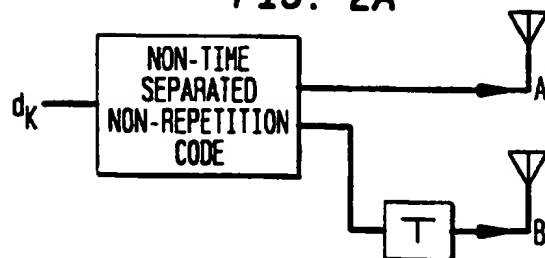


FIG. 2B

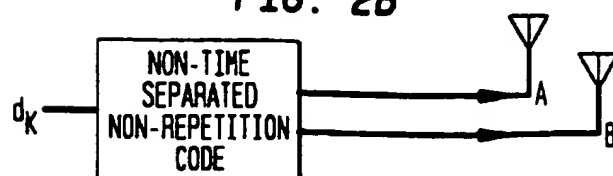




FIG. 3

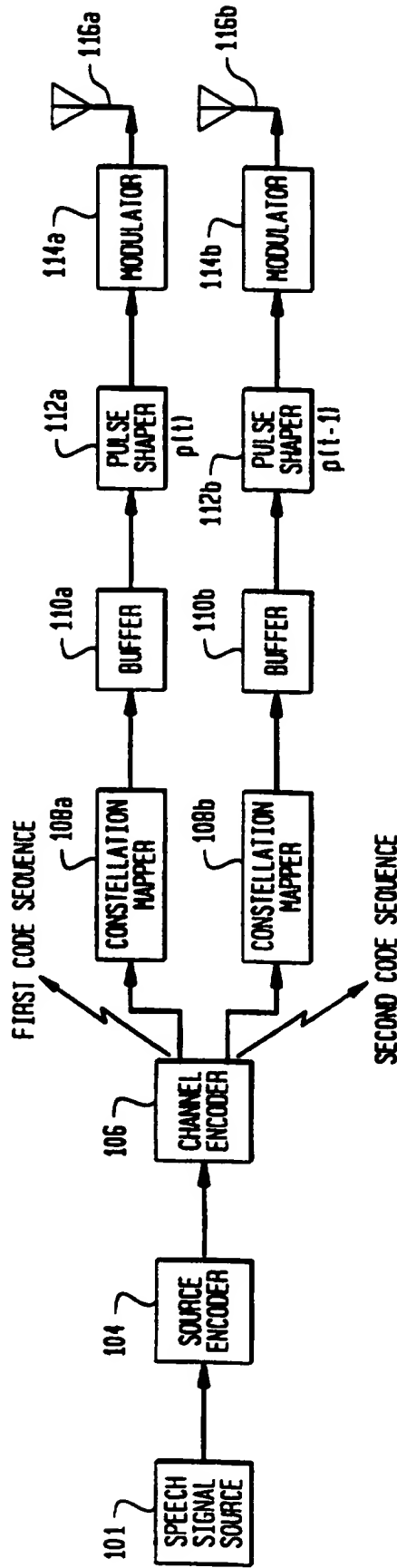
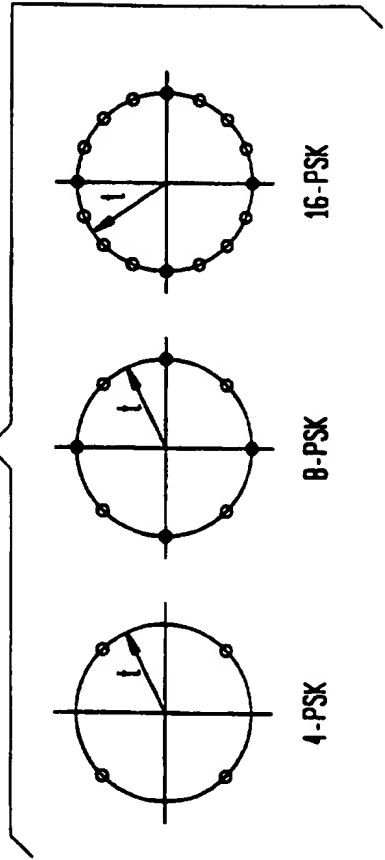


FIG. 4



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FIG. 5

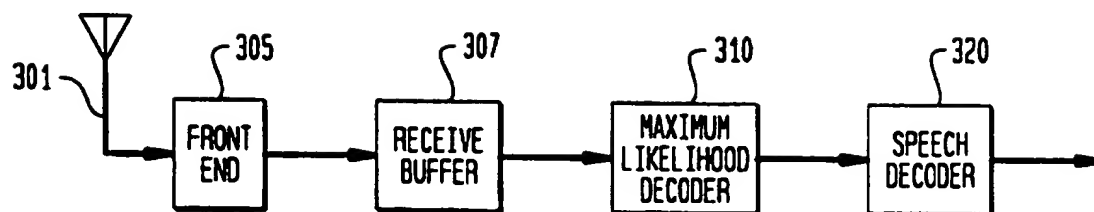
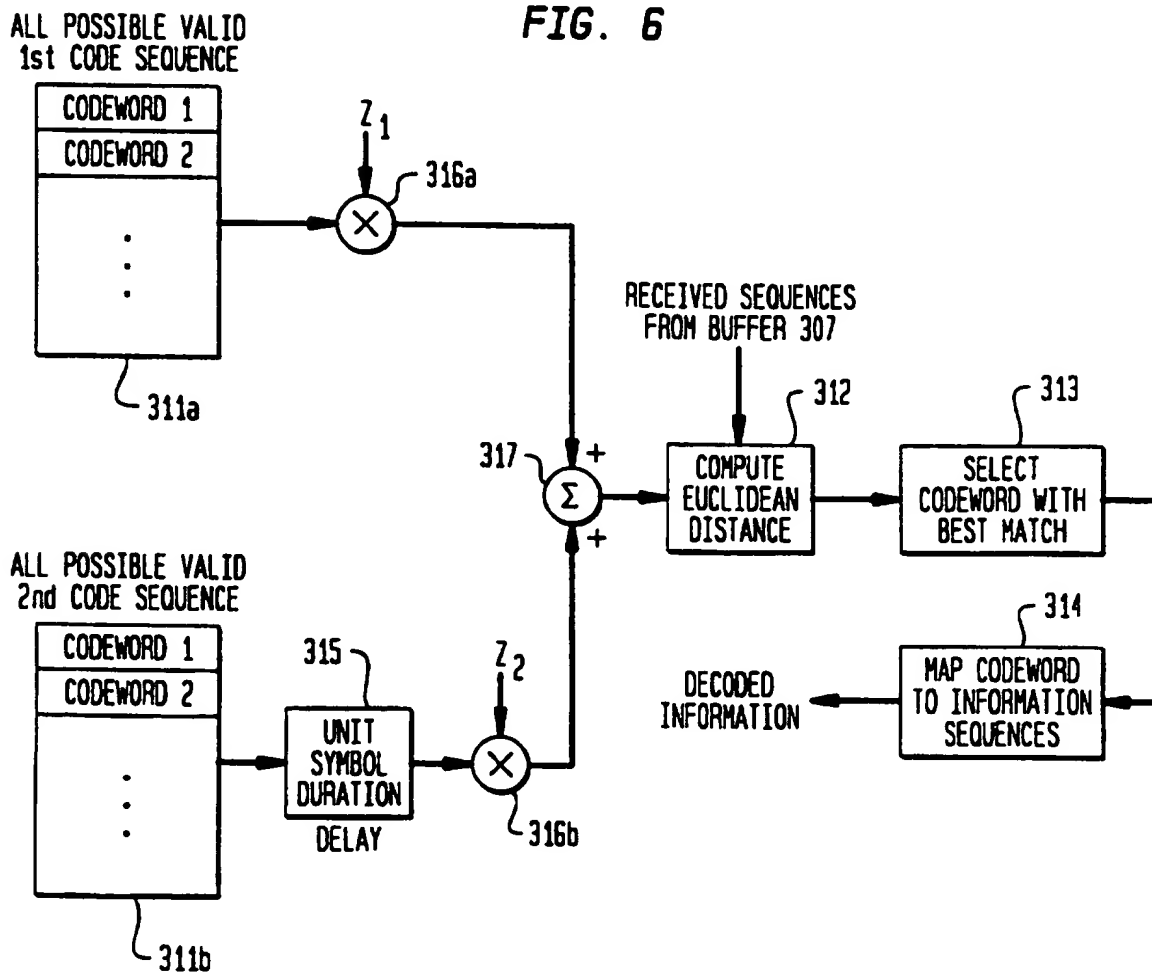


FIG. 6



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FIG. 7

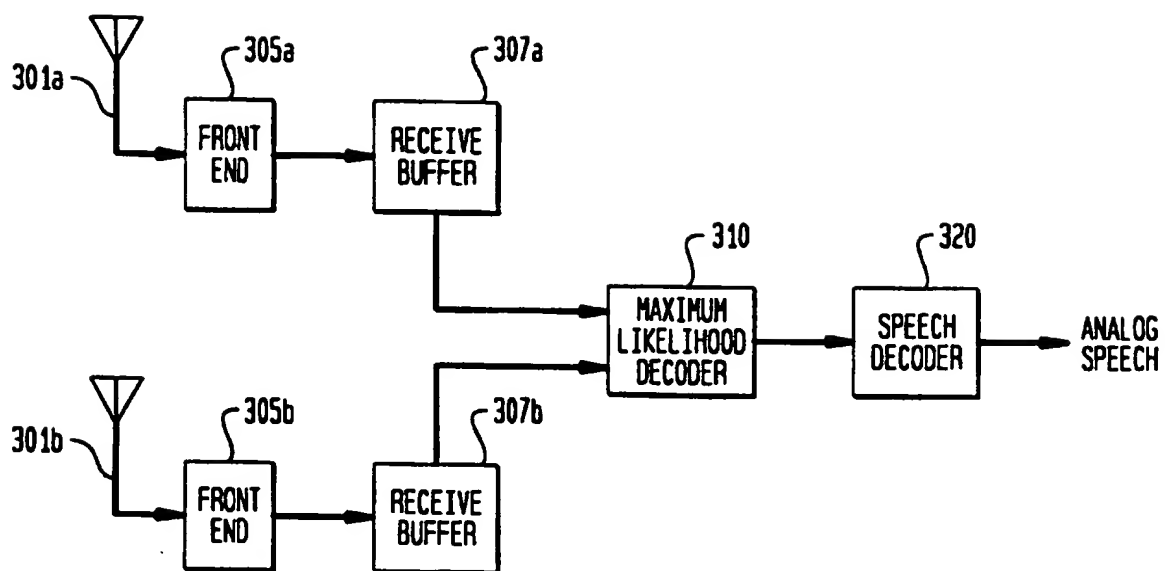


FIG. 8

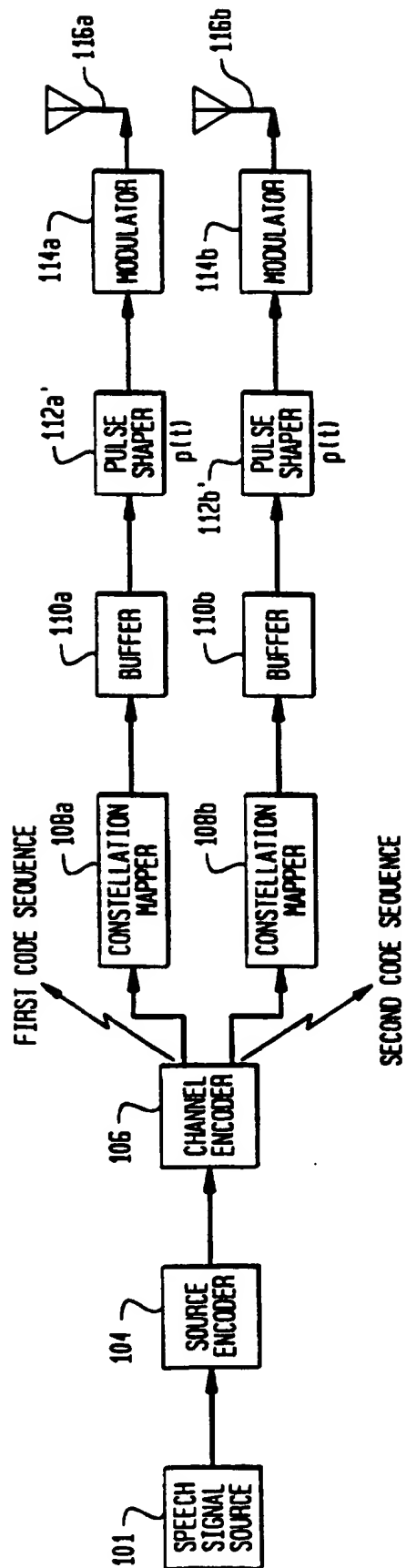
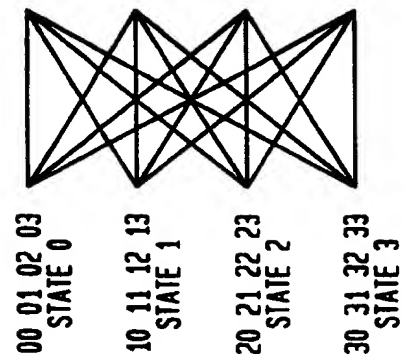
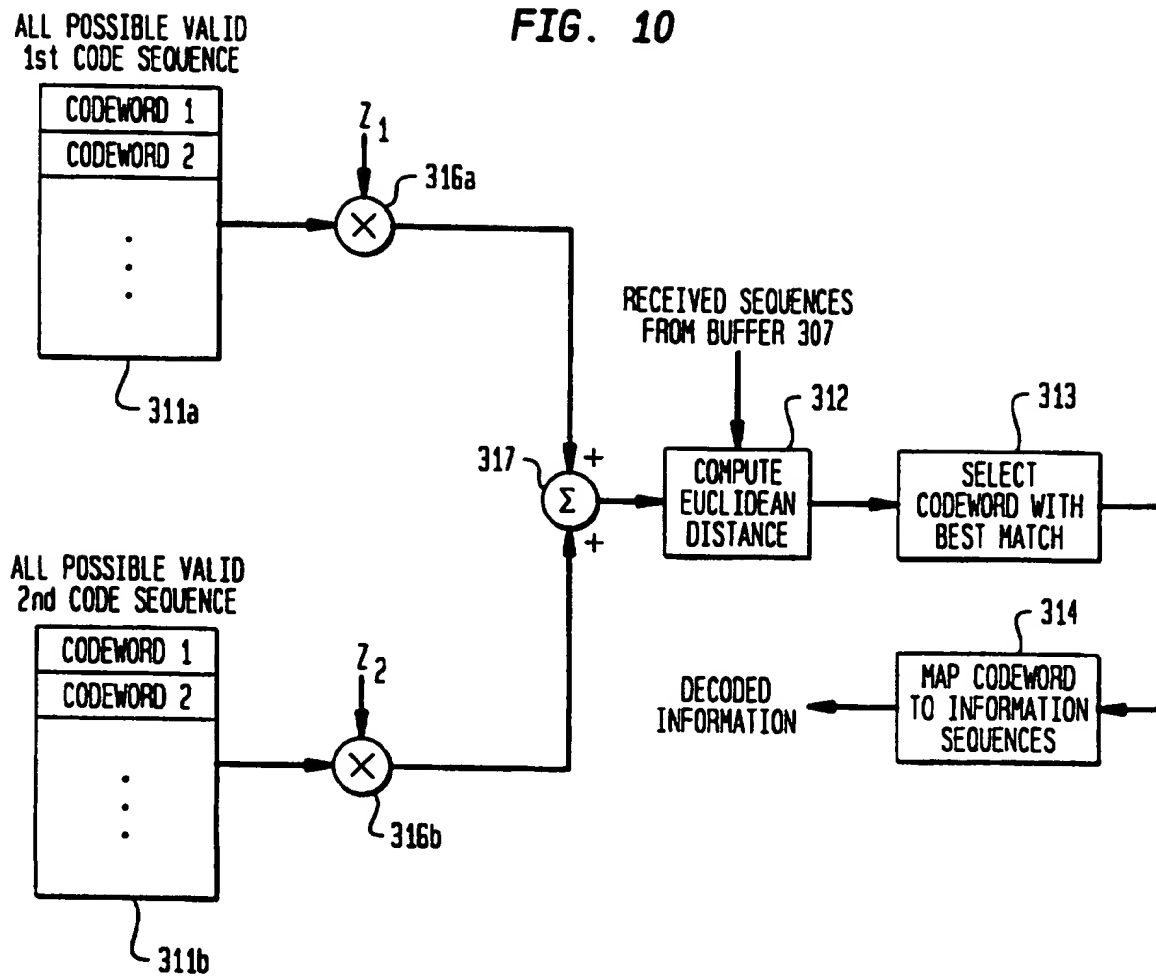


FIG. 9



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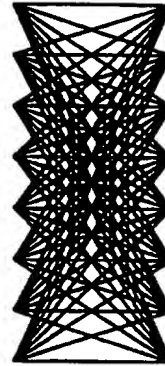
FIG. 10



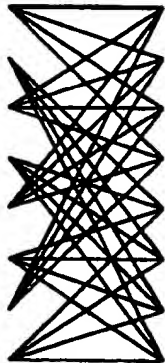
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**FIG. 11**

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 70. 71. 72. 73. 74. 75. 76. 77  
 40. 41. 42. 43. 44. 45. 46. 47  
 10. 11. 12. 13. 14. 15. 16. 17  
 60. 61. 62. 63. 64. 65. 66. 67  
 30. 31. 32. 33. 34. 35. 36. 37

**FIG. 12**

00. 01. 02. 03  
 10. 11. 12. 13  
 20. 21. 22. 23  
 30. 31. 32. 33  
 22. 23. 20. 21  
 32. 33. 30. 31  
 02. 03. 00. 01  
 12. 13. 10. 11

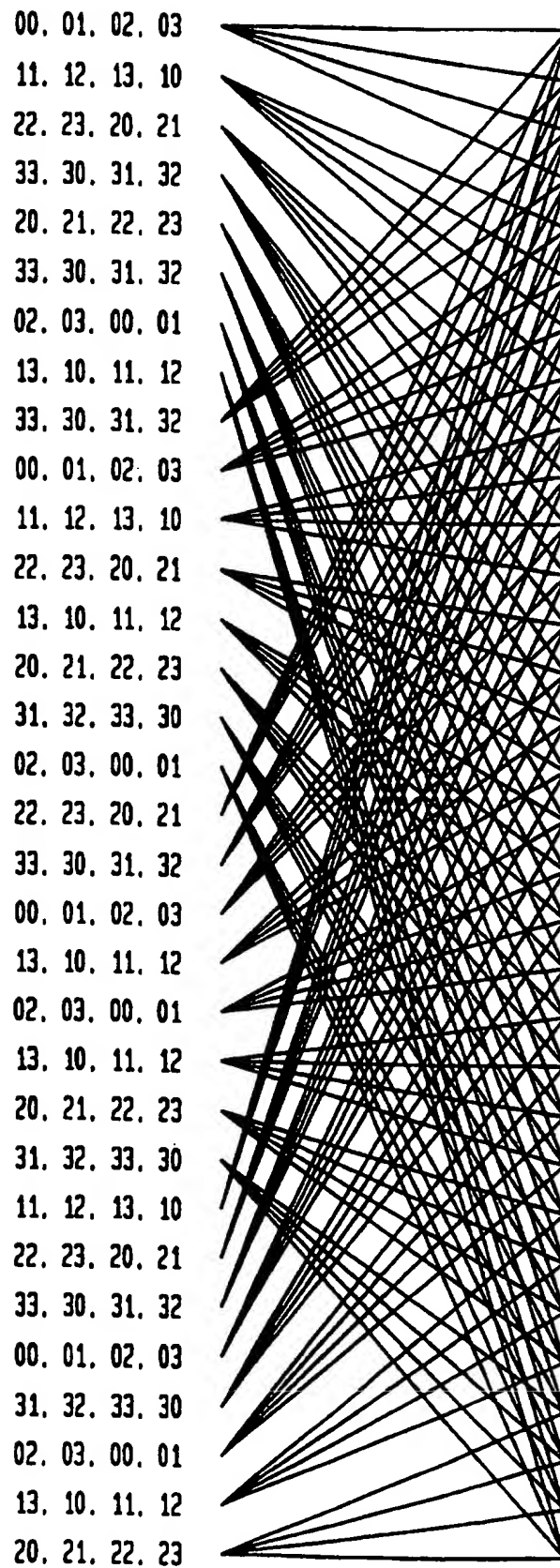


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 12. 13. 10. 11  
 02. 03. 00. 01  
 10. 11. 12. 13  
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 10. 11. 12. 13

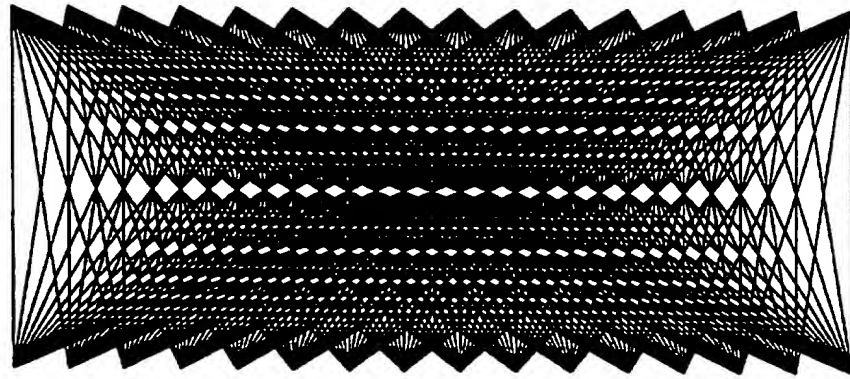


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FIG. 13



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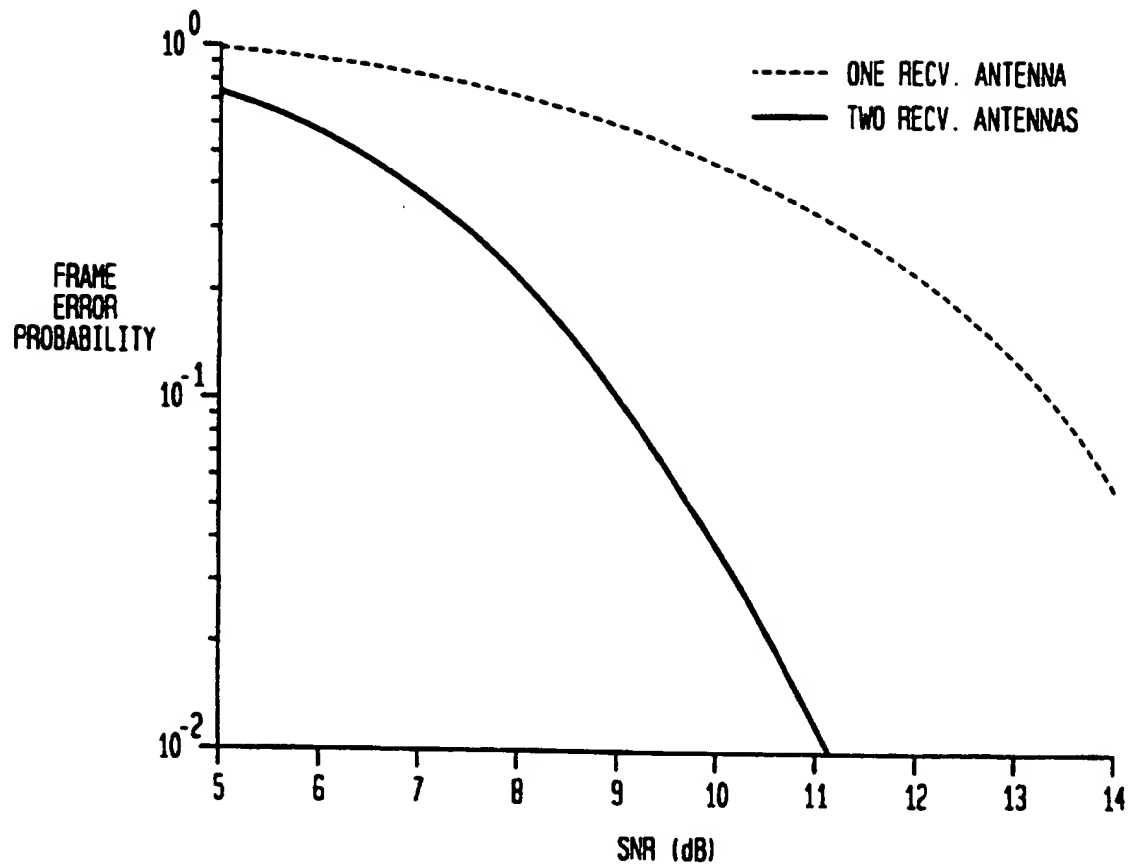
**FIG. 14**

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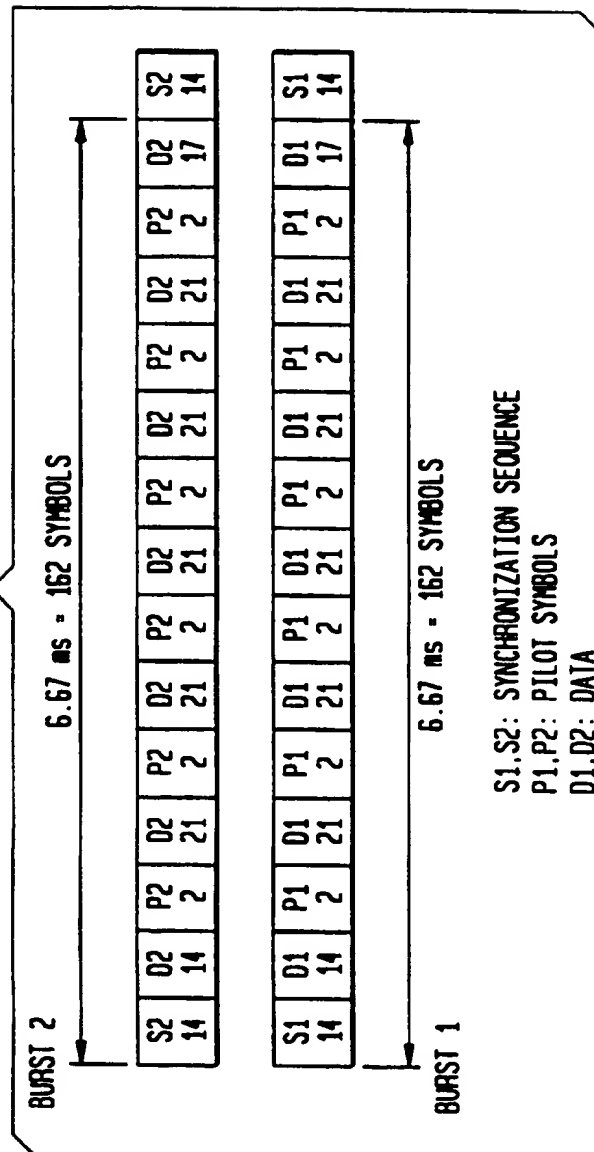


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FIG. 15



**FIG. 16**



**FIG. 17**

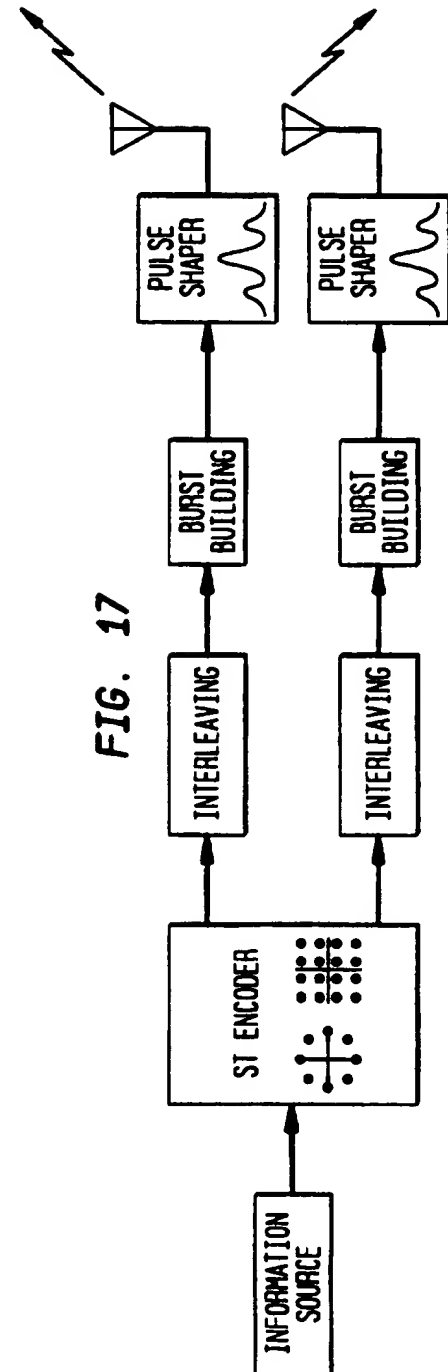


FIG. 18

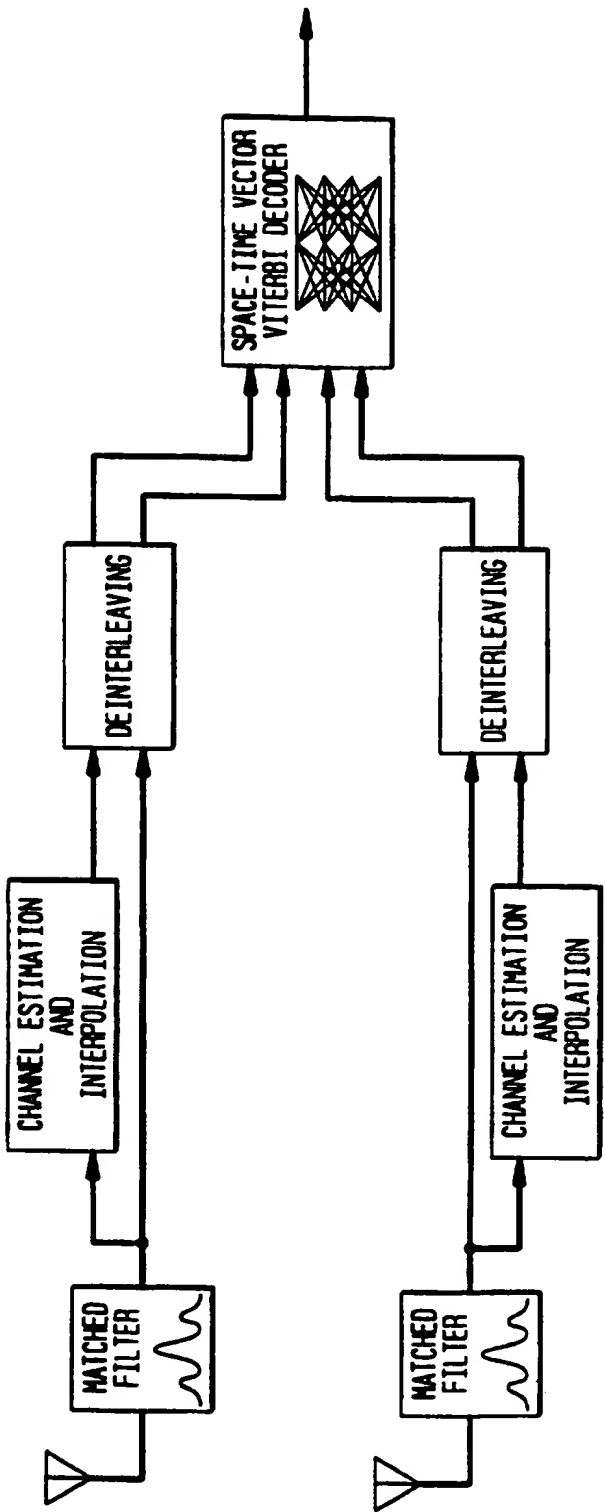
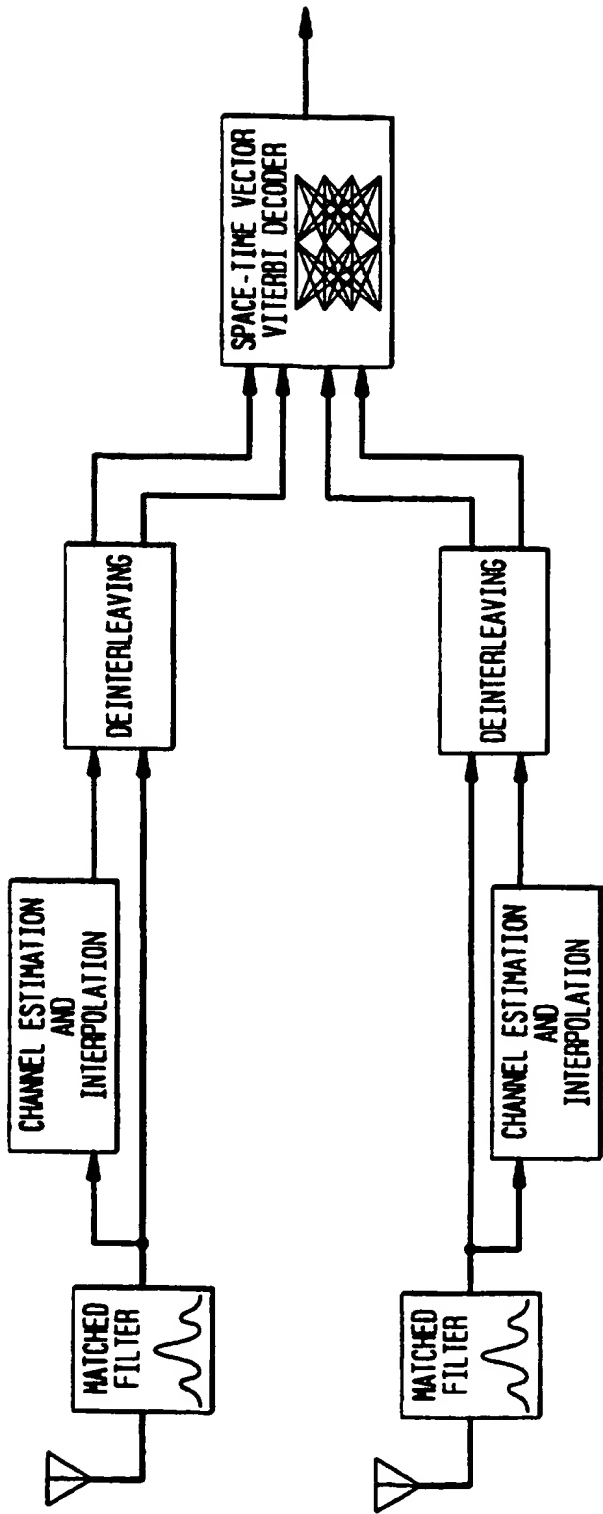


FIG. 18



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FIG. 19

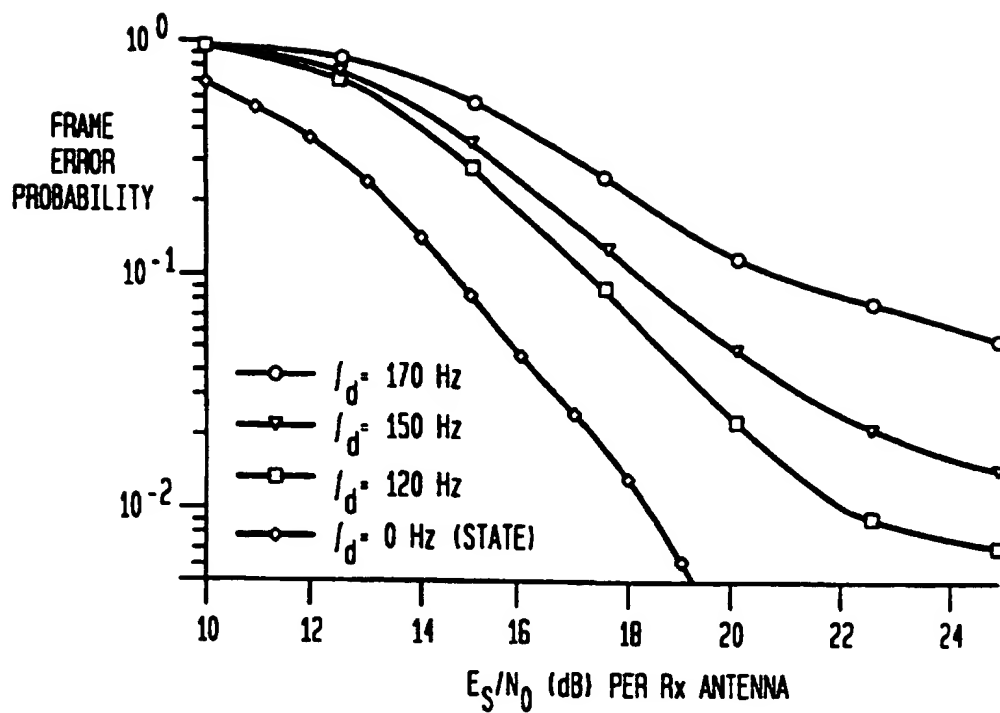
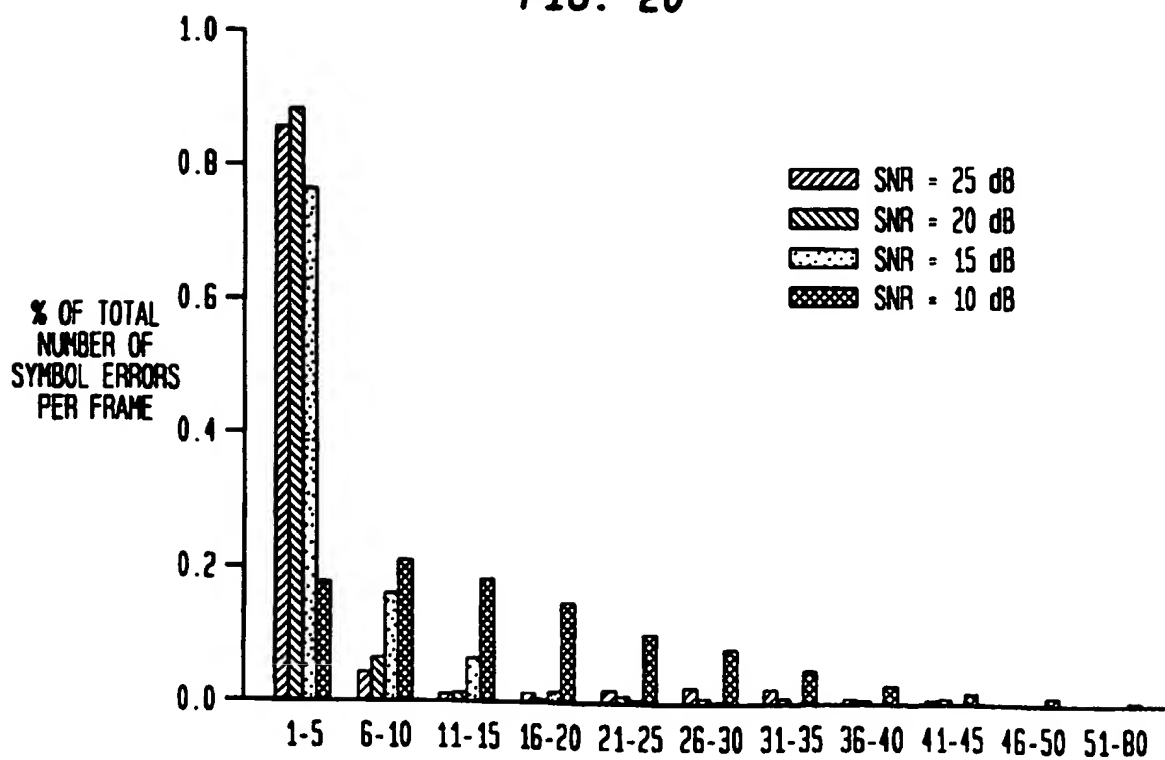
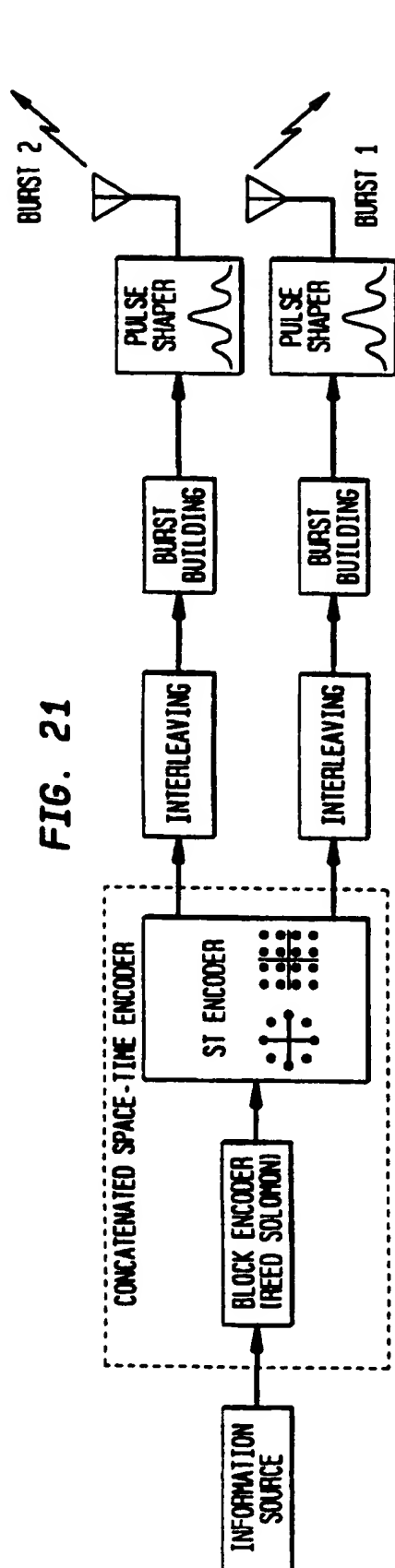
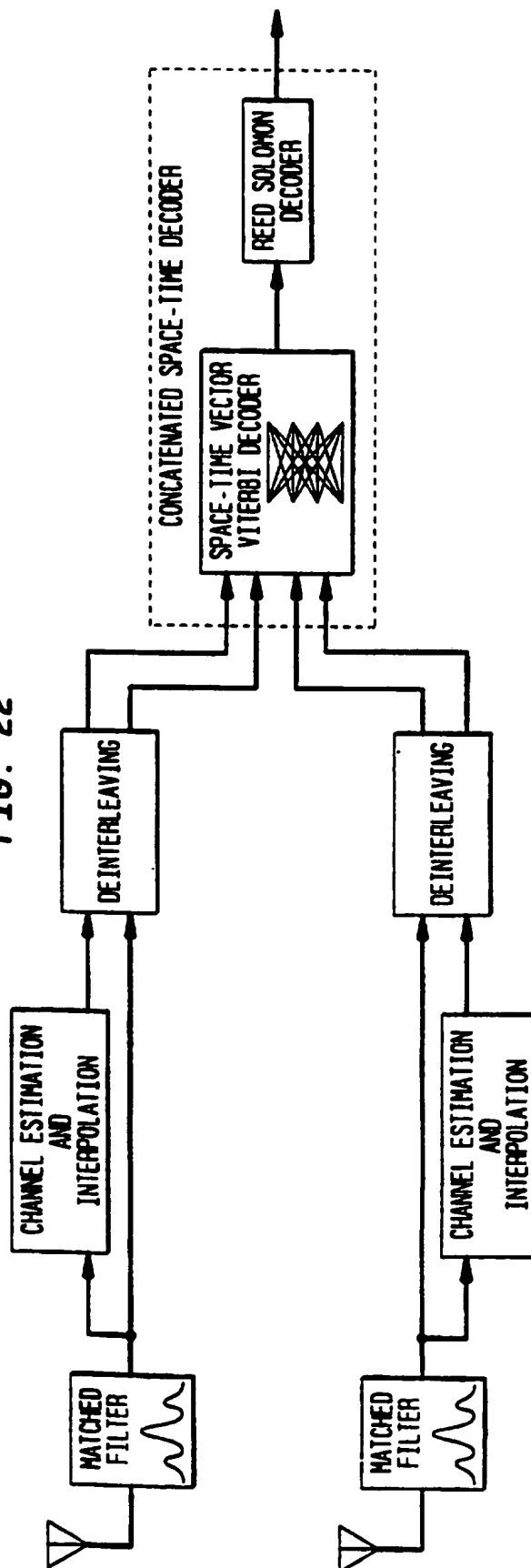


FIG. 20



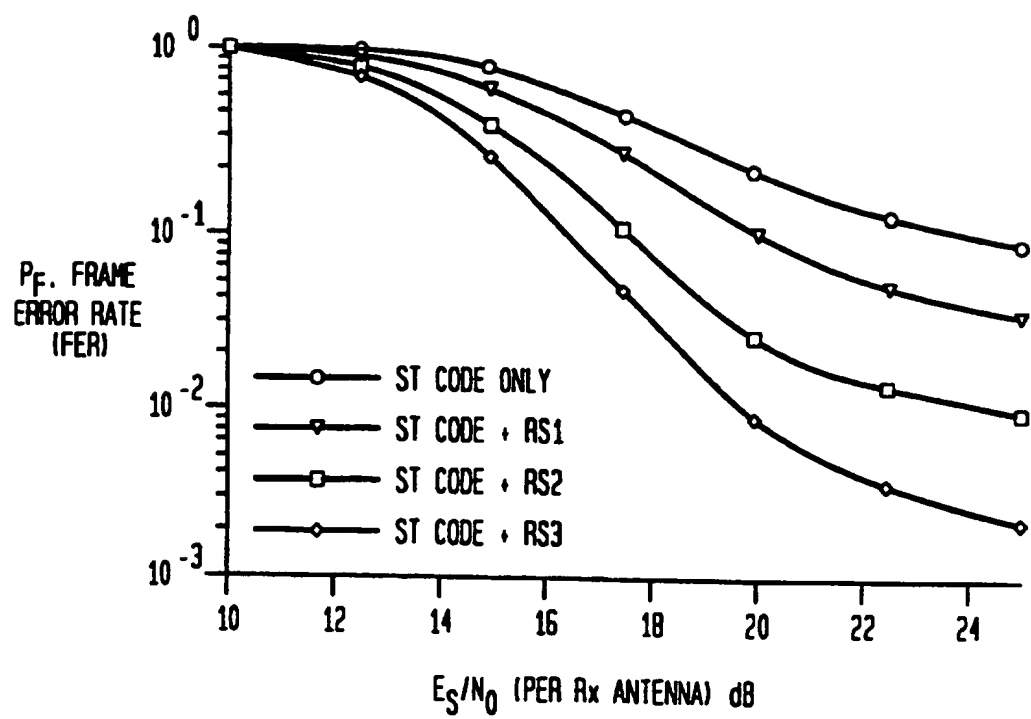


**FIG. 22**



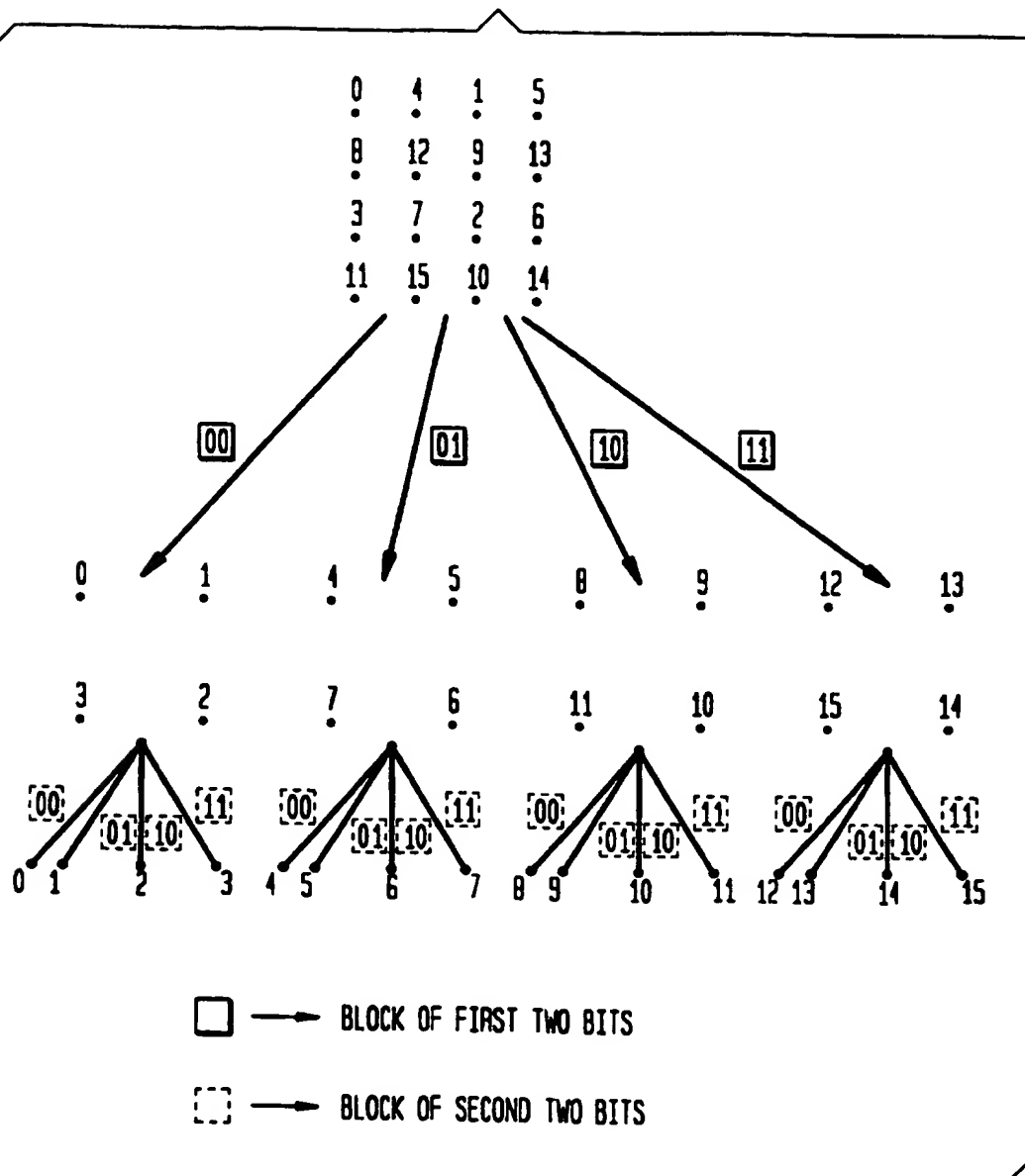
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FIG. 23



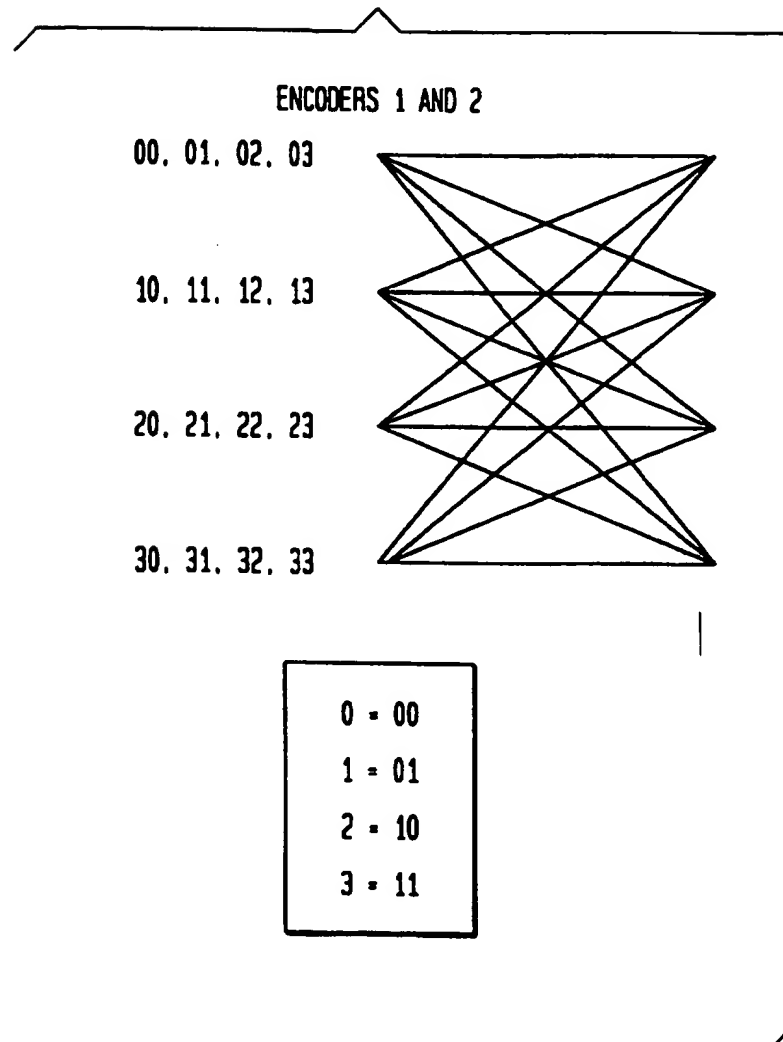
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FIG. 24

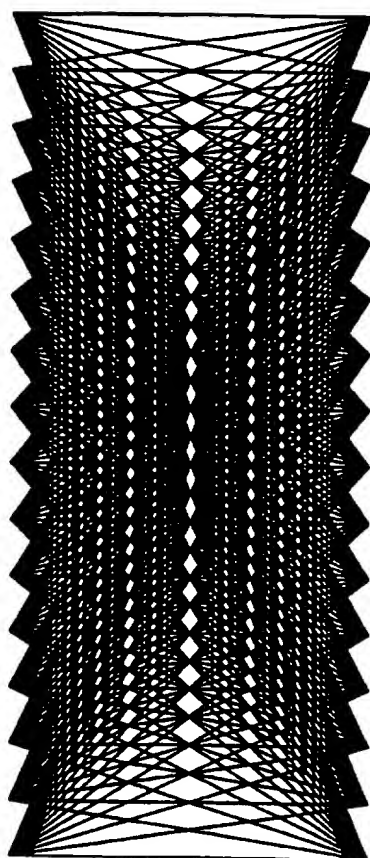




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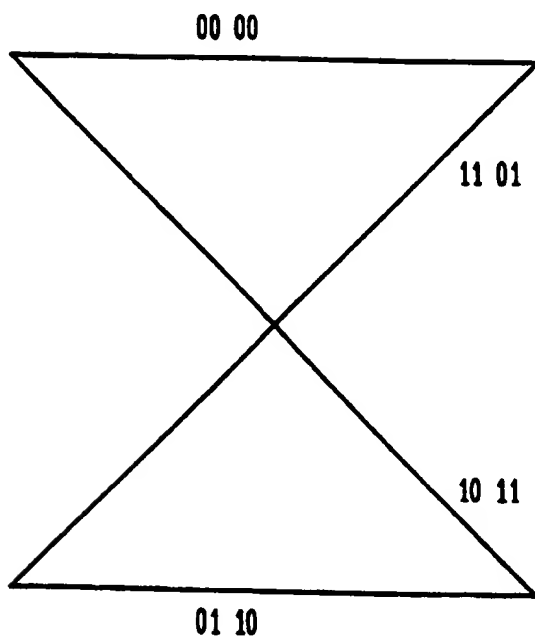
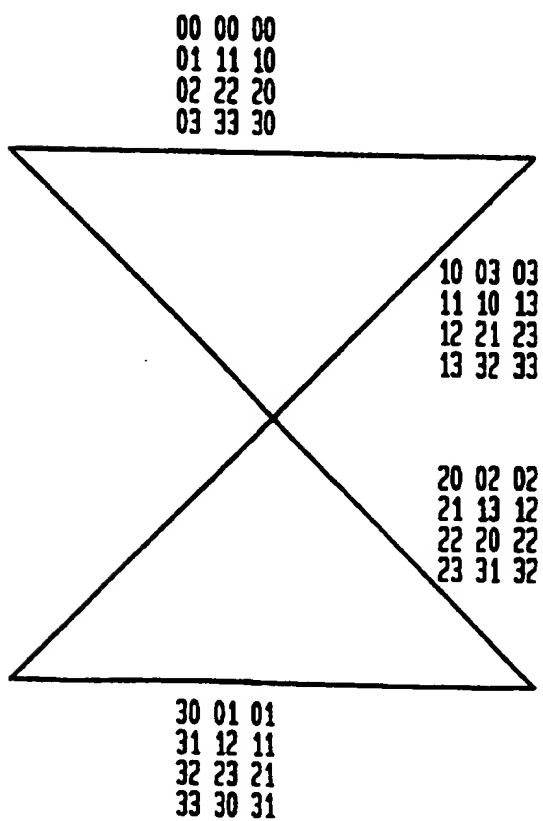
**FIG. 25**

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**FIG. 26**

00. 01. 02. ... . 015  
10. 11. 12. ... . 115  
20. 21. 22. ... . 215  
30. 31. 32. ... . 315  
40. 41. 42. ... . 415  
50. 51. 52. ... . 515  
60. 61. 62. ... . 615  
70. 71. 72. ... . 715  
80. 81. 82. ... . 815  
90. 91. 92. ... . 915  
100. 101. ... . 1015  
110. 111. ... . 1115  
120. 121. ... . 1215  
130. 131. ... . 1315  
140. 141. ... . 1415  
150. 151. ... . 1515

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**FIG. 27A****FIG. 27B**

## INTERNATIONAL SEARCH REPORT

International application No.  
PCT/US97/07010

**A. CLASSIFICATION OF SUBJECT MATTER**

IPC(6) : H04L 23/02

US CL : 375/299,265

According to International Patent Classification (IPC) or to both national classification and IPC

**B. FIELDS SEARCHED**

Minimum documentation searched (classification system followed by classification symbols)

U.S. : 375/299,265,260-264,267,340,341

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practicable, search terms used)

APS

search terms: diversity, trellis, branches.

**C. DOCUMENTS CONSIDERED TO BE RELEVANT**

| Category* | Citation of document, with indication, where appropriate, of the relevant passages          | Relevant to claim No.                              |
|-----------|---|--|
| X         | US 5,305,353 A (WEERACKODY) 19 April 1994, see Figs. 2,3 and col.5, line 6 - col.6, line 61 | 1-3, 7-10, 15-23, 43, 44, 47, 50-53, 56-60, 64, 65 |
| ----      |   | -----  |
| Y         |   | 11, 24, 42, 46, 54, 63                             |
| X         | US 5,396,518 A (HOW) 07 March 1995, see col.2, lines 25-45.                                 | 42-44, 63-65                                       |
| X         | US 5,418,798 A (WEI) 23 May 1995, see Fig. 1.   | 16-21, 59, 60                                      |
| Y         | US 3,633,107 A (BRADY) 04 January 1972, see cols. 1, 2.                                     | 11, 24, 42, 46, 54, 63                             |

☐ Further documents are listed in the continuation of Box C. ☐ See patent family annex.

|   |      |  |
|---|------|--|
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| * E* earlier document published on or after the international filing date   | * Y  | document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art |
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| * O* document referring to an oral disclosure, use, exhibition or other means   |      |  |
| * P* document published prior to the international filing date but later than the priority date claimed   |      |  |

Date of the actual completion of the international search

23 JUNE 1997

Date of mailing of the international search report

28 AUG 1997

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